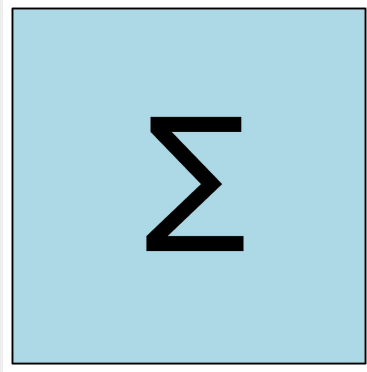


Practical zk with Summa

Napoli - 25/5/2023



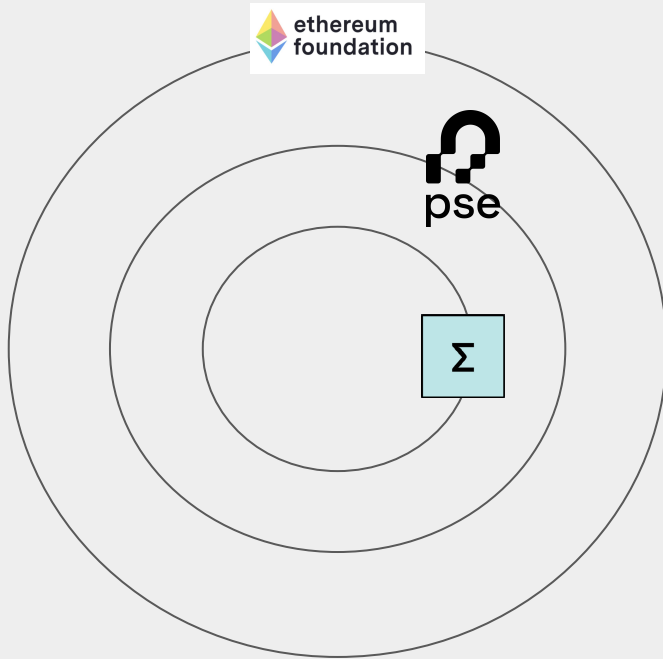
enrico.eth

Me

- Zero knowledge circuit engineer at Summa
- Previously at Polygon Hermez

Summa


- zk Proof of Solvency for CEXes



Goal : Understand zk in practice

#1 Context on ZK

#2 Context on something practical

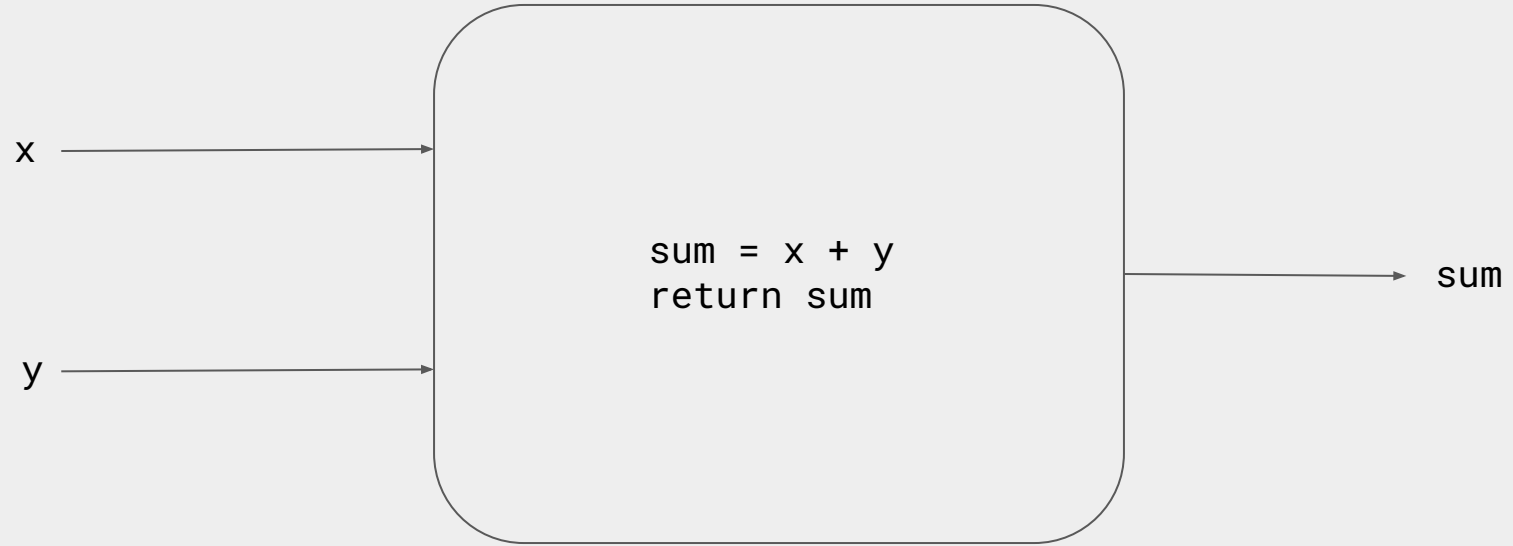


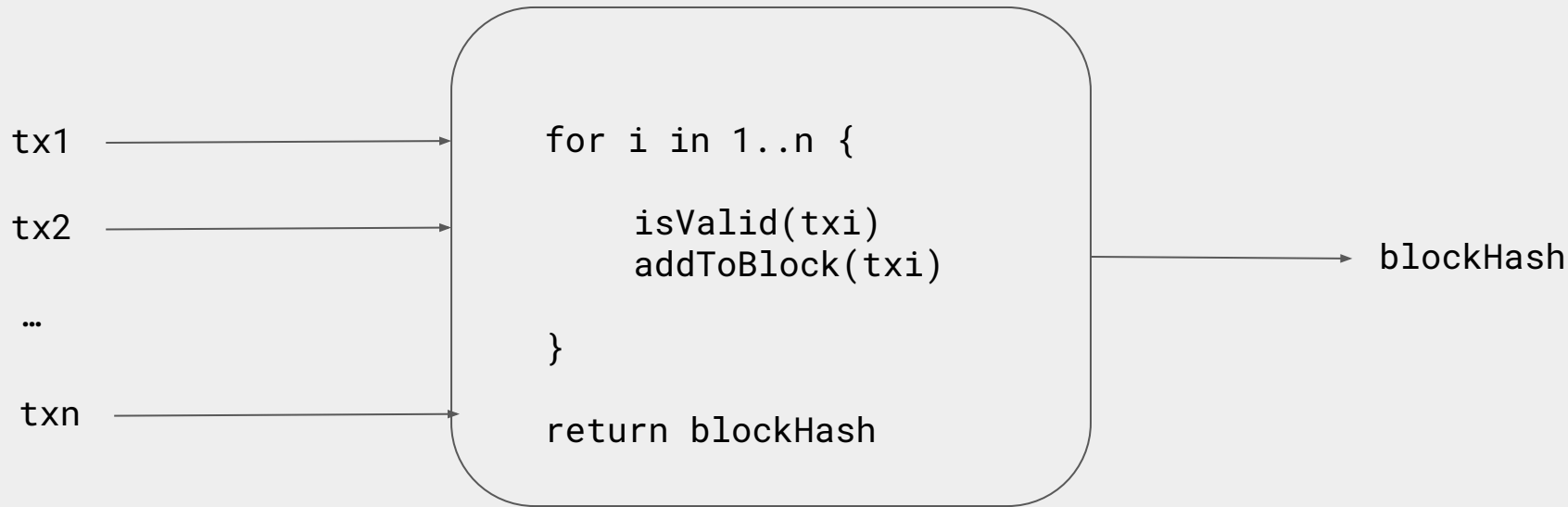
#3 Apply zk to something practical

#1 Context on ZK

Computation = set of rules

```
// Function that computes the sum of two numbers
function computeSum(x: number, y: number): number {
  // Rule 1: Compute the sum of 'x' and 'y' and return the result.
  const sum: number = x + y;
  return sum;
}
```

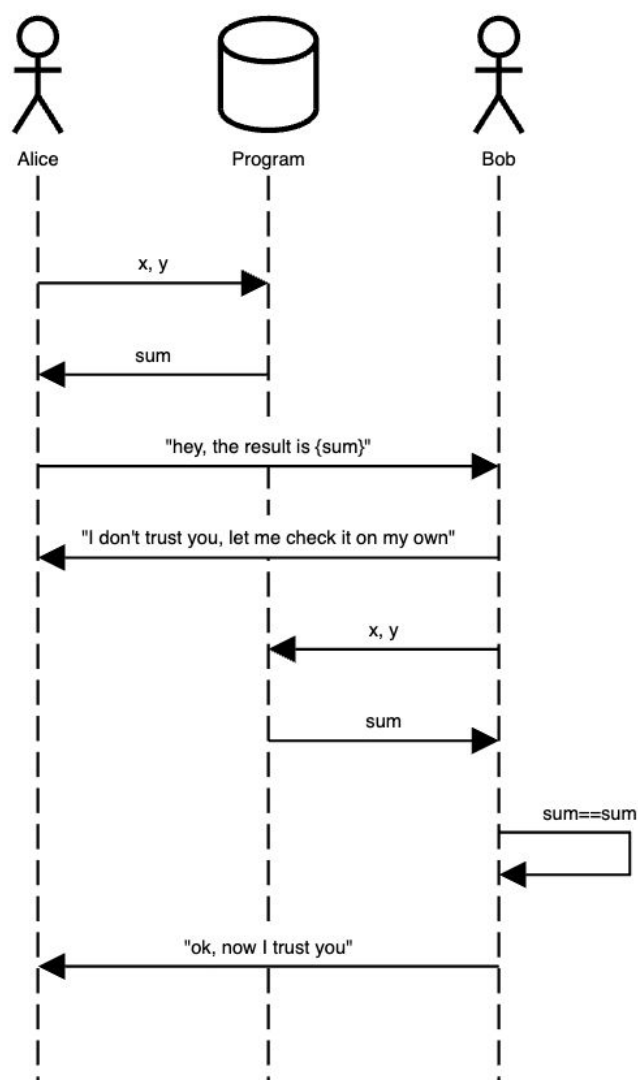




Computational Integrity Guarantee

Given a computation whose rules are known by everyone, a **prover** wants to prove that the output is the result of running the computation on certain inputs

Traditional Prover Verifier dynamic

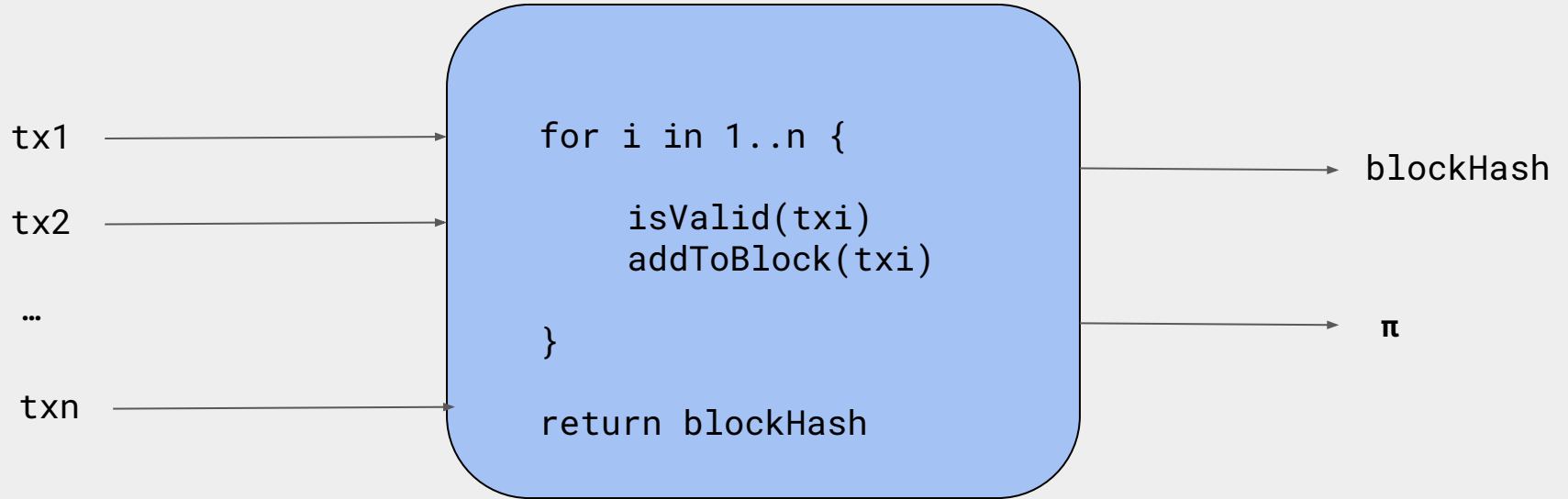


How to achieve Computational Integrity Guarantee?

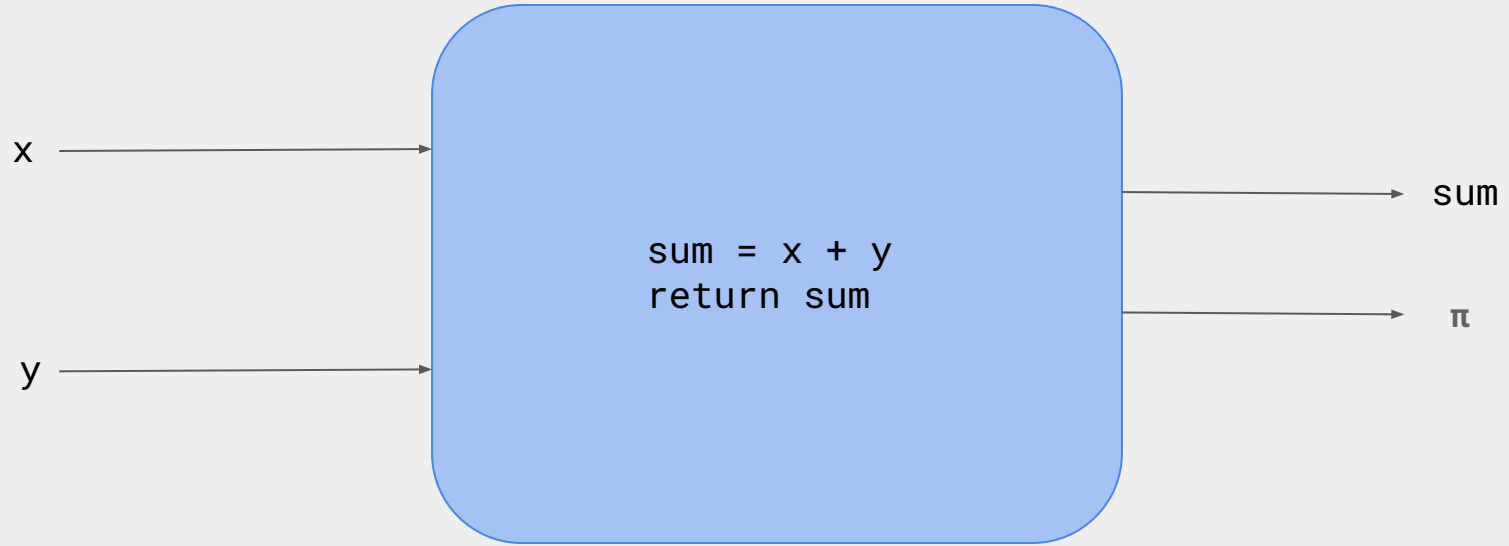
- Verifier needs to rerun the computation with the same input and check that the output matches
- Issue #1: verification time is linear to the computation
- Issue #2: everything is public!

Enters ZK

zk Circuit



zk Circuit

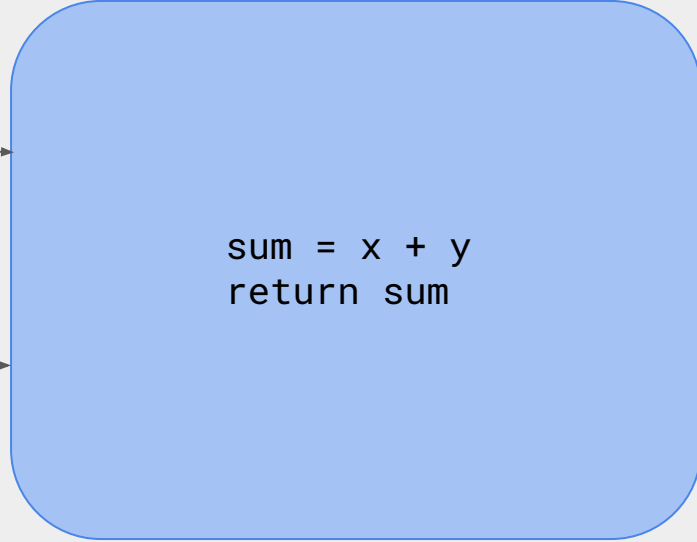


zk Circuit

penultimate_level_left_hash



y

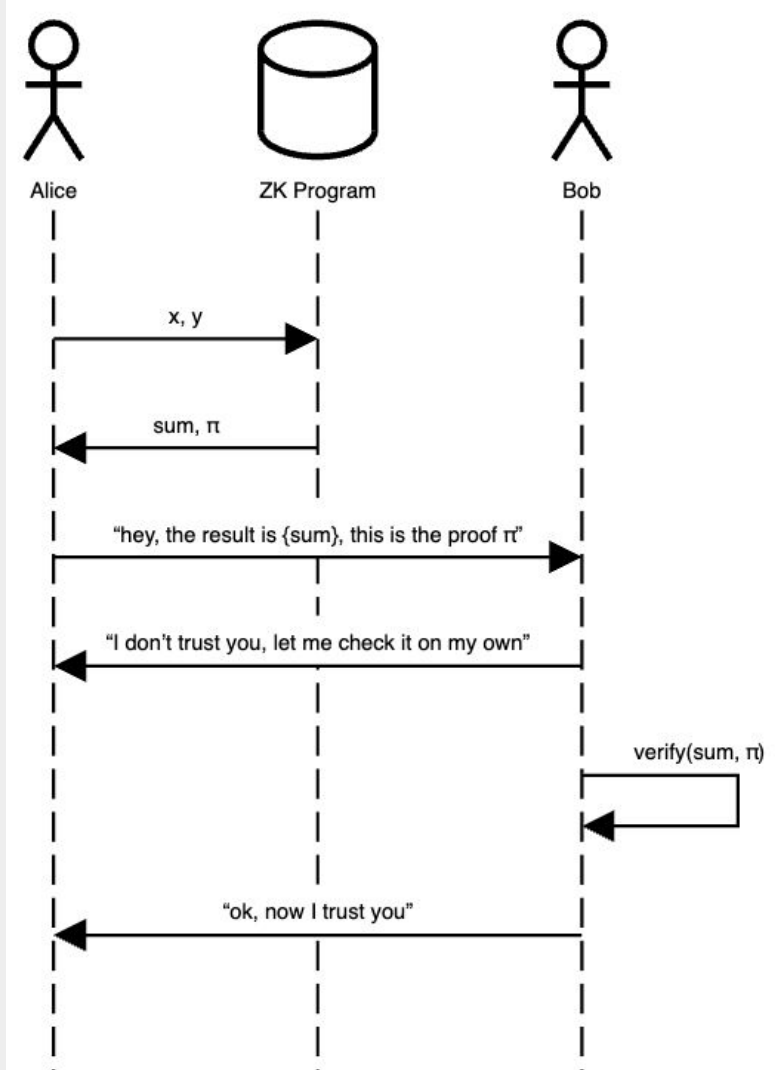


sum



π

ZK Prover Verifier dynamic

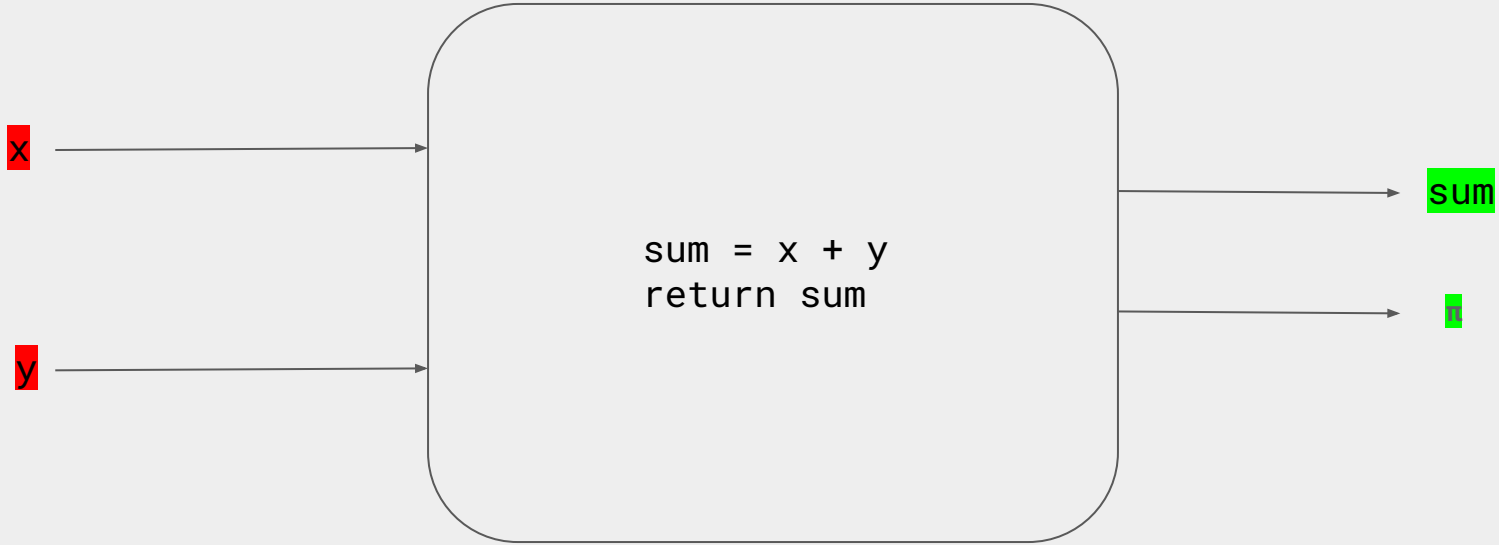


How to achieve Computational Integrity Guarantee?

- Verifier needs to run a verification algorithm on the proof π
- Solved Issue #1: verification time is constant no matter the time it took to run the computation
- Solved Issue #2: the prover can selectively decide what to keep private and what to keep public

How to achieve Computational Integrity Guarantee?

- Verifier needs to run a verification algorithm on the proof π
- Solved Issue #1: verification time is constant no matter the time it took to run the computation
- Solved Issue #2: the prover can selectively decide what to keep private and what to keep public
- New issue #1: Generating a proof for a computation is way slower than just running the computation
- New issue #2: Writing zk program is not as easy as writing a normal program



#2 Context on Something Practical

Proof of Solvency

- Cryptographic proof that a CEX is solvent at a specific moment in time

Proof of Solvency

- Cryptographic proof that a CEX is solvent at a specific moment in time



Assets \geq Liabilities

LIABILITIES

- Deposits of the users
- Denominated in ETH, BTC, USDC ...
- Do not live on-chain, live in the CEX's DB

ASSETS

- Cryptographic assets (ETH, BTC, USDC...) controlled by the CEX
- Live on-chain
- Should map (at least) 1:1 the deposits of the users

LIABILITIES

- Deposits of the users
- Denominated in ETH, BTC, USDC ...
- Do not live on-chain, live in the CEX's DB

Proof Of Solvency

- Cryptographic proof that a CEX is solvent at a specific moment in time



Assets \geq Liabilities



Users are confident
that they can withdraw
at any time

#3 Apply ZK to something practical

Summa: ZK Proof of Solvency

auditor-based proof of solvency



everything is
ok!



auditorless proof of solvency (naive approach)



everything is
ok!

auditorless proof of solvency (ZK approach)

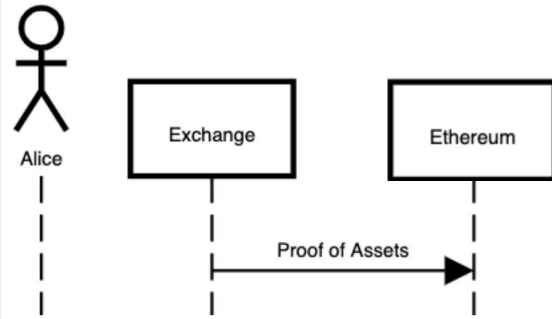


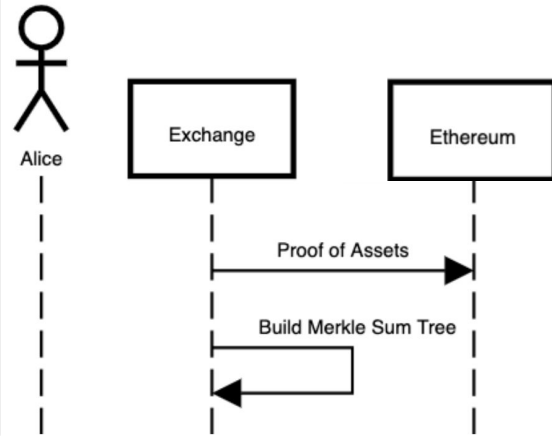
Π

everything is
ok!

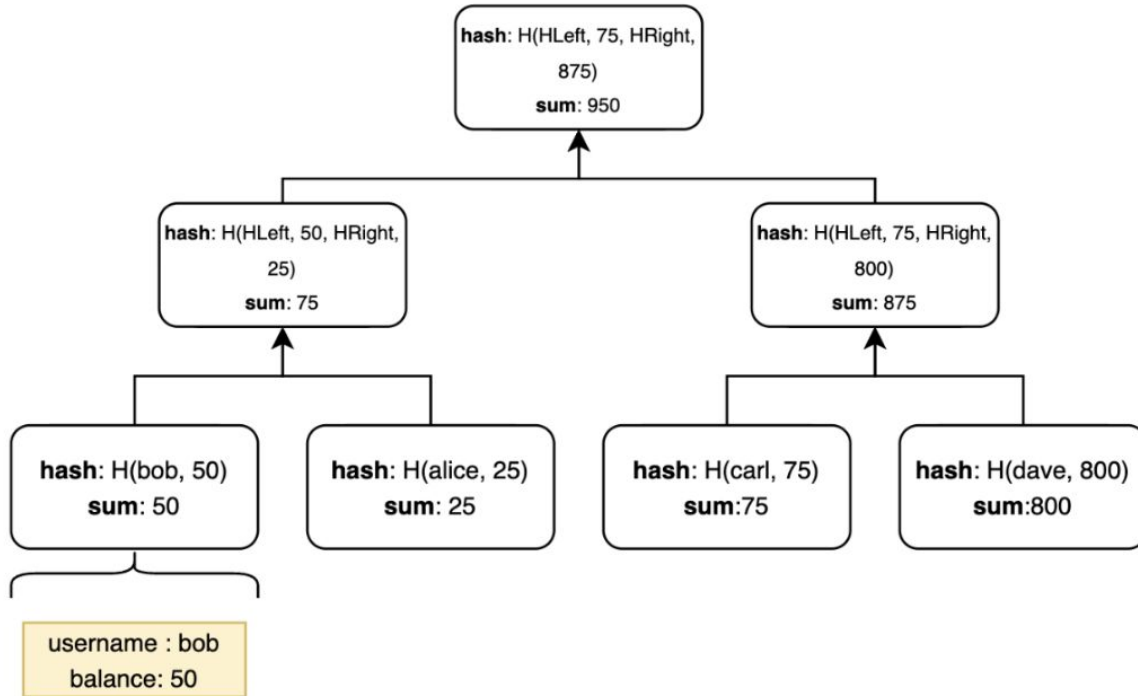


How?

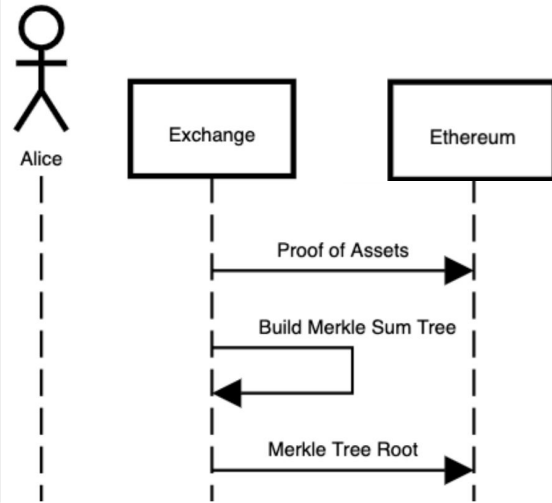


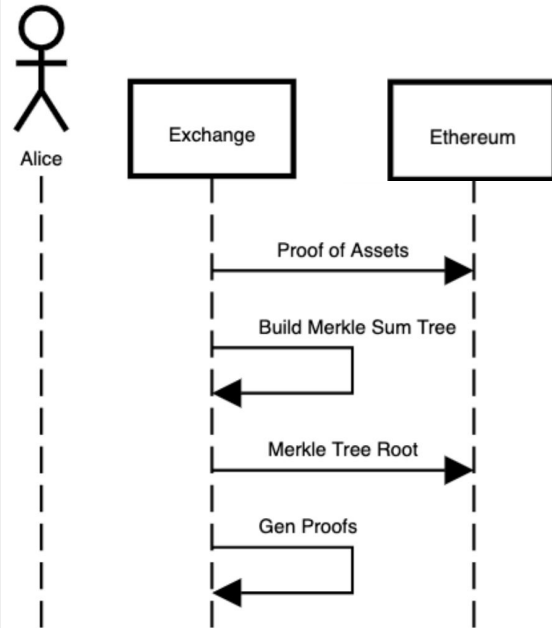


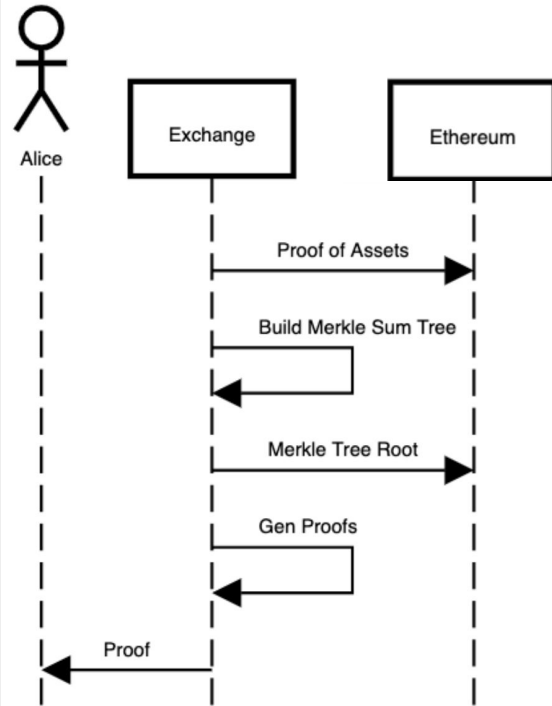
Merkle Sum Tree



- The entries are the users' data (= liabilities)
- Lives off-chain
- Only the root-hash gets published on-chain

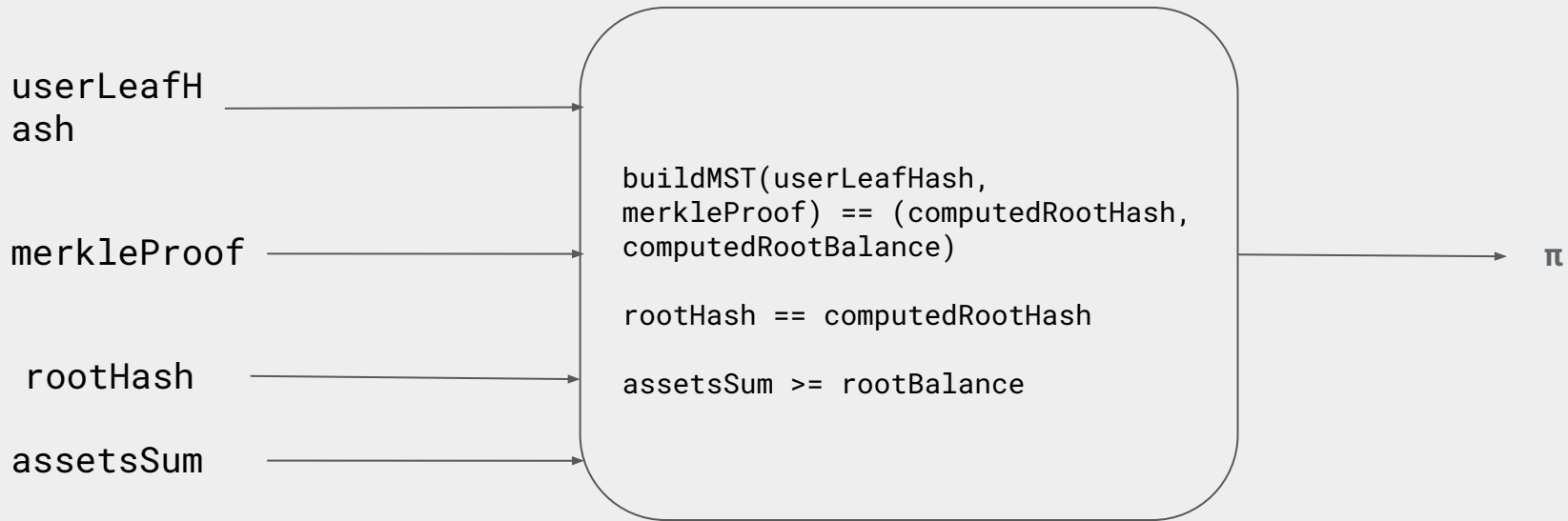






Zk Proofs - Program Rules

- Rule#1: The user (identified by its username) is included in the Merkle Sum Tree with the correct balance
- Rule#2: The hash of the Merkle Sum Tree matches the one committed on chain
- Rule#3: The sum of liabilities is Less Than the assets of the exchange (as committed in step 1)
- Rule#4: No sum overflow happened in the merkle sum tree computation



Zk Proofs - secrecy

- Other users information such as their balances and usernames
- Total number of users
- Total amount of liabilities

userLeafHash

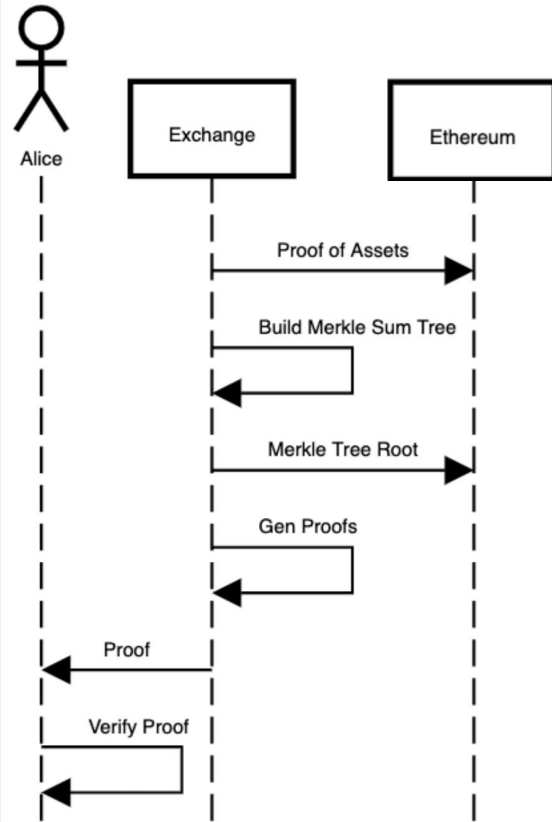
merkleProof

rootHash

assetsSum

```
buildMST(userLeafHash,  
merkleProof) == (computedRootHash,  
computedRootBalance)  
  
rootHash == computedRootHash  
  
assetsSum >= rootBalance
```

!!



Proof Verification

$F(\pi, \text{username}, \text{balance}, \text{assetsSum}, \text{rootHash}) =$
yes/no

Conclusions

→ How to think of zk apps: the mental model

- ◆ Given a computation which rules are known by everyone, a prover wants to prove that the output is the result of running the computation on certain inputs, without revealing (part of) the input of such computation

→ How to build zk apps

- ◆ Building zk apps means writing circuits. Circom is the best tool to get started

Thank you!
me on github

