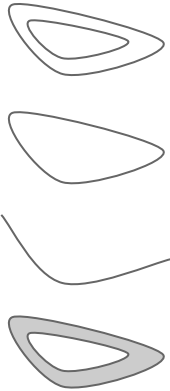


# Tessellated GPU Path Rendering

Stephen White, GPU-MON

# What's a path?

- path: unordered set of contours
- contour: ordered set of segments
- segment: line, quadratic, cubic, conic
- may have holes, may be self-intersecting
- winding rule to specify how holes are filled
- Web: SVG, <canvas>, rounded rect, fonts



# Why do we even need to do this?

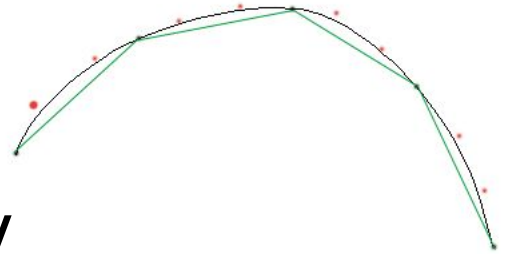
- GPUs render very fast, but only triangles
- must convert from path to triangles

# Several approaches

- rasterize in software, upload to texture
- stencil-and-cover
- distance fields
- tessellation
- Loop-Blinn (plus tessellation)

# How to handle curved segments?

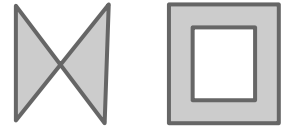
- linearize them (subdivision)
- contours now piecewise-linear
- use screenspace threshold
- problem now: tessellate arbitrary polygonal contours into triangles



# Animals in the polygon zoo

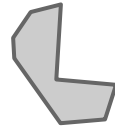
- **Complex**

- self-intersections, holes



- **Simple**

- no self-intersections, no holes, concave



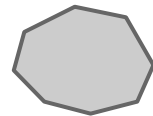
- **Monotone**

- monotonically incr; all pts one side, concave



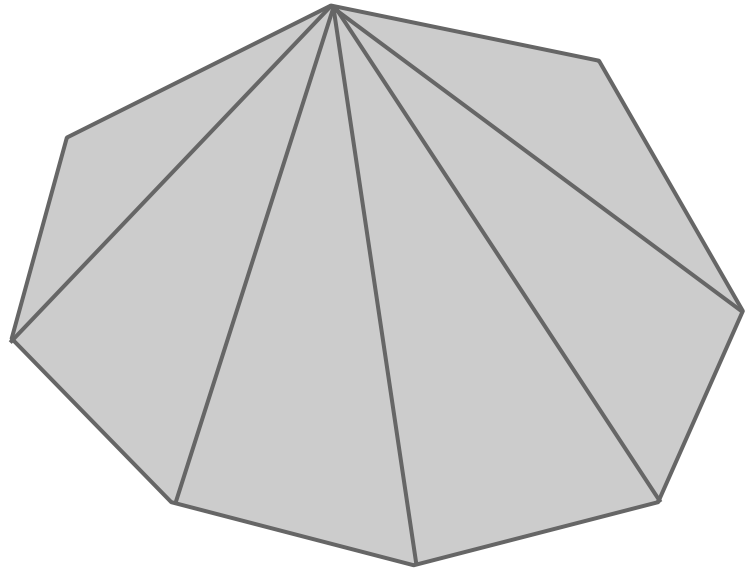
- **Convex**

- all angles  $\leq 180^\circ$



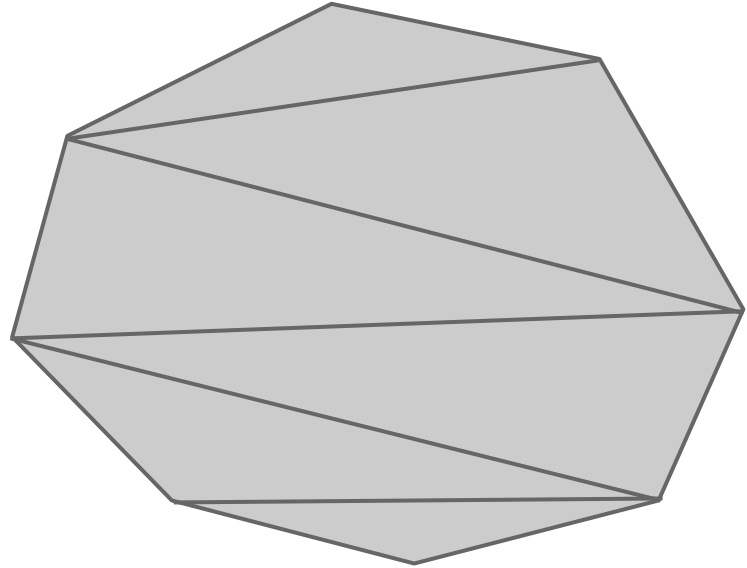
# Convex polygons

First approach: fan



# Convex polygons

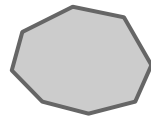
Second approach: Alternate sides



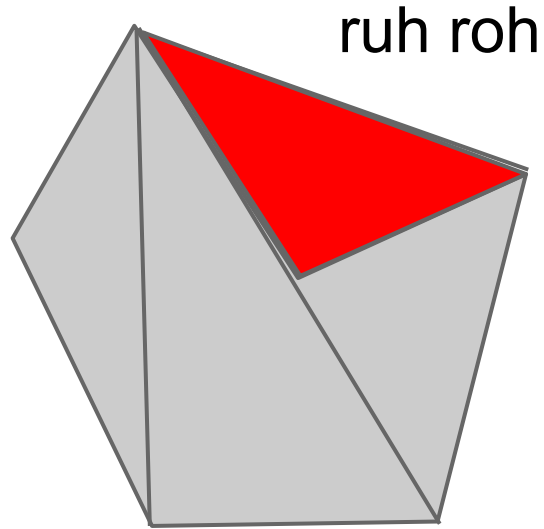


# Animals in the polygon zoo

- Complex
  - self-intersections, holes
- Simple
  - no self-intersections, no holes, concave
- Monotone
  - monotonically incr; all pts one side, concave
- **Convex**
  - all angles  $\leq 180^\circ$



# Concave: doesn't work



# Concave: doesn't work

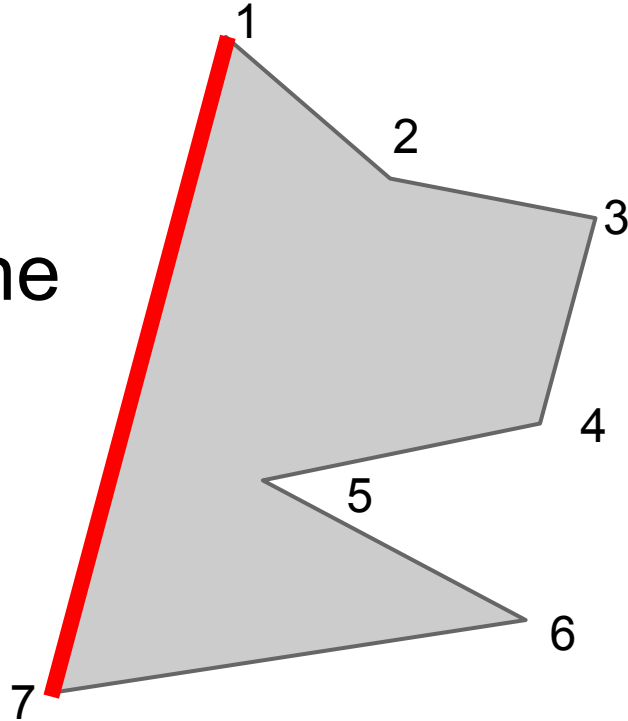
“If at first you don't succeed ... try something easier.”

- Alfred E. Neumann



# Monotone polygon

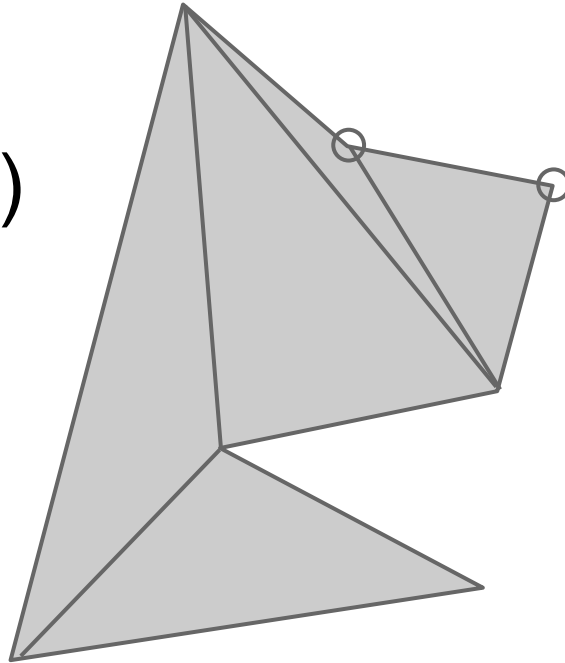
All points lie to one side of the minY->maxY line



Points are monotonically increasing in Y

# Triangulate monotones: ear clipping

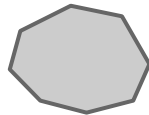
- 1) find the first convex vertex ( $O(n)$ )
- 2) clip its ear ( $O(1)$ )
- 3) find convex neighbour
- 4) clip its ear
- 5) repeat



backtrack at most  
one step  $\Rightarrow O(n)$

# Animals in the polygon zoo

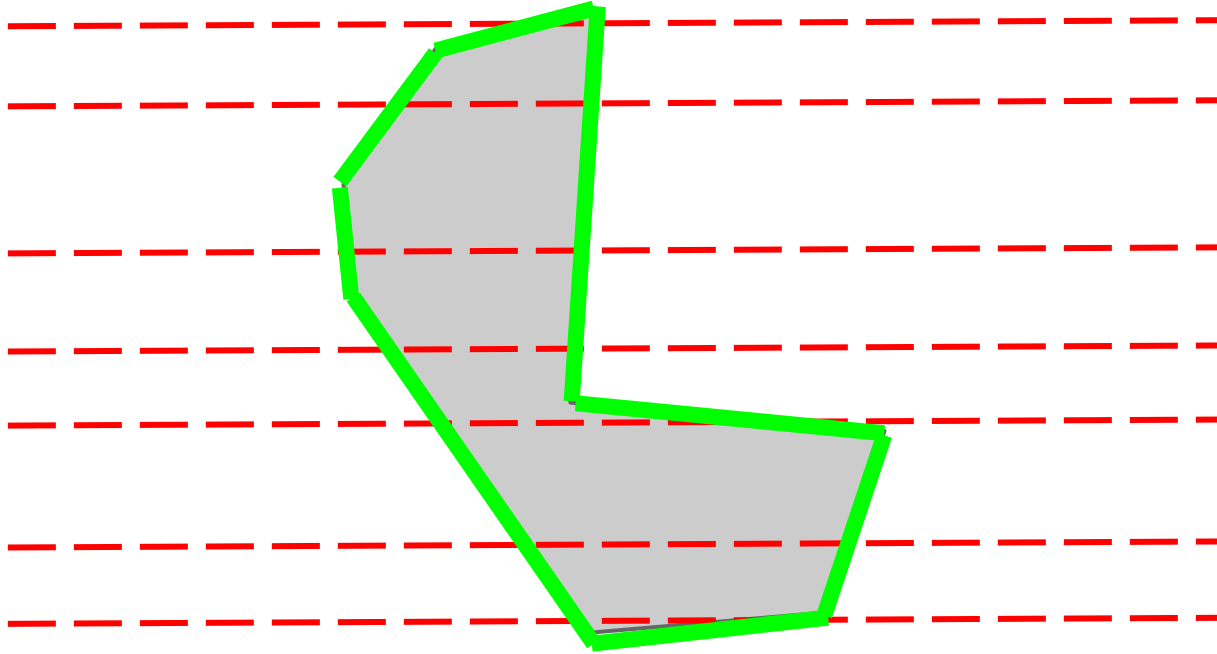
- Complex
  - self-intersections, holes
- Simple
  - no self-intersections, no holes, concave
- Monotone
  - monotonically incr; all pts one side, concave
- Convex
  - all angles  $\leq 180^\circ$



# Sweep Line Algorithms

- Sort vertices in  $Y$  ( $O(N \lg N)$ )
- Sweep line passes from top to bottom
- Active Edge List: top vertex seen, bottom not
- Two vertices on the same sweep line?
  - WLOG: rotate slightly in  $Z$
- $y_1 == y_2 \Rightarrow$  Secondary compare in  $X$
- “enclosing edges” of a vertex

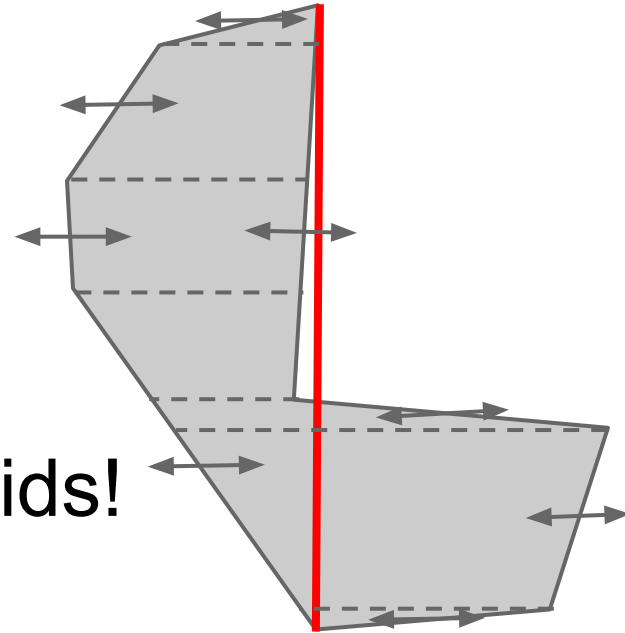
# Active edge list: example





# Lil' bit harder now: simple polygon

- Obviously, not monotonic
- Trapezoidal decomposition: turn simple into monotones
- I say, forget about the trapezoids!
- Which monotone lies left/right of each edge?



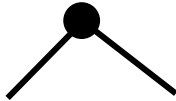
# Lil' bit harder now: simple polygon

Categorize verts in the zoo

All degree 2, so three species:



type 1:  
one edge above,  
one edge below



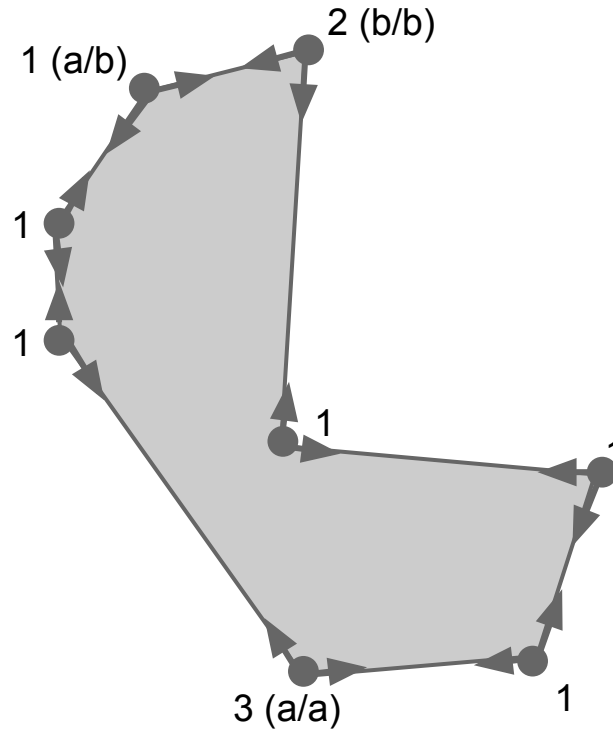
type 2:  
both edges below



type 3:  
both edges above

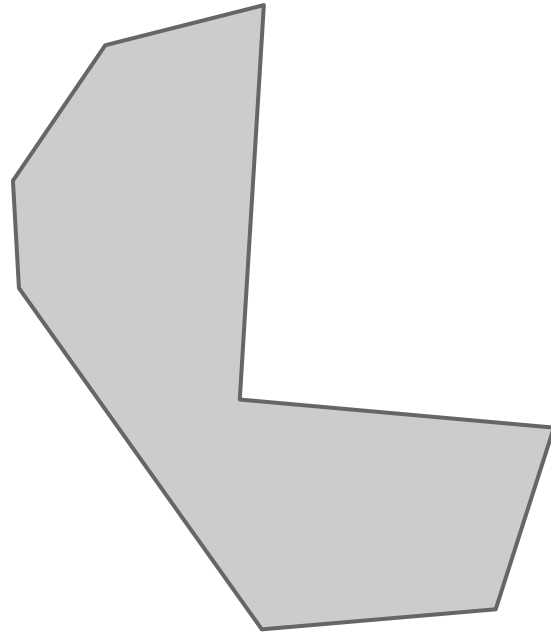
# Lil' bit harder now: simple polygon

- vertex types



# Lil' bit harder now: simple polygon

- Sort vertices in  $Y$   
(secondarily in  $X$ )
- Each edge stores left  
mono, right mono

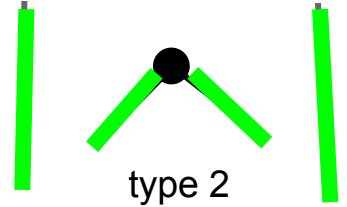


# Simple to monotone: da rules

- type 1:
  - AEL: remove incoming, insert outgoing
- type 2:
  - AEL: find enclosing edges (if any), and insert new edges in between
- type 3:
  - AEL: remove incoming edges



type 1



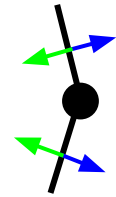
type 2



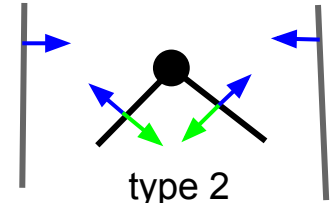
type 3

# Simple to monotone: da rules

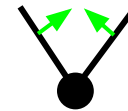
- type 1:
  - copy L/R monotones from incoming
- type 2:
  - copy enclosing monotone (if any) to exterior edges
  - open new monotone below
- type 3:
  - close monotone above



type 1



type 2

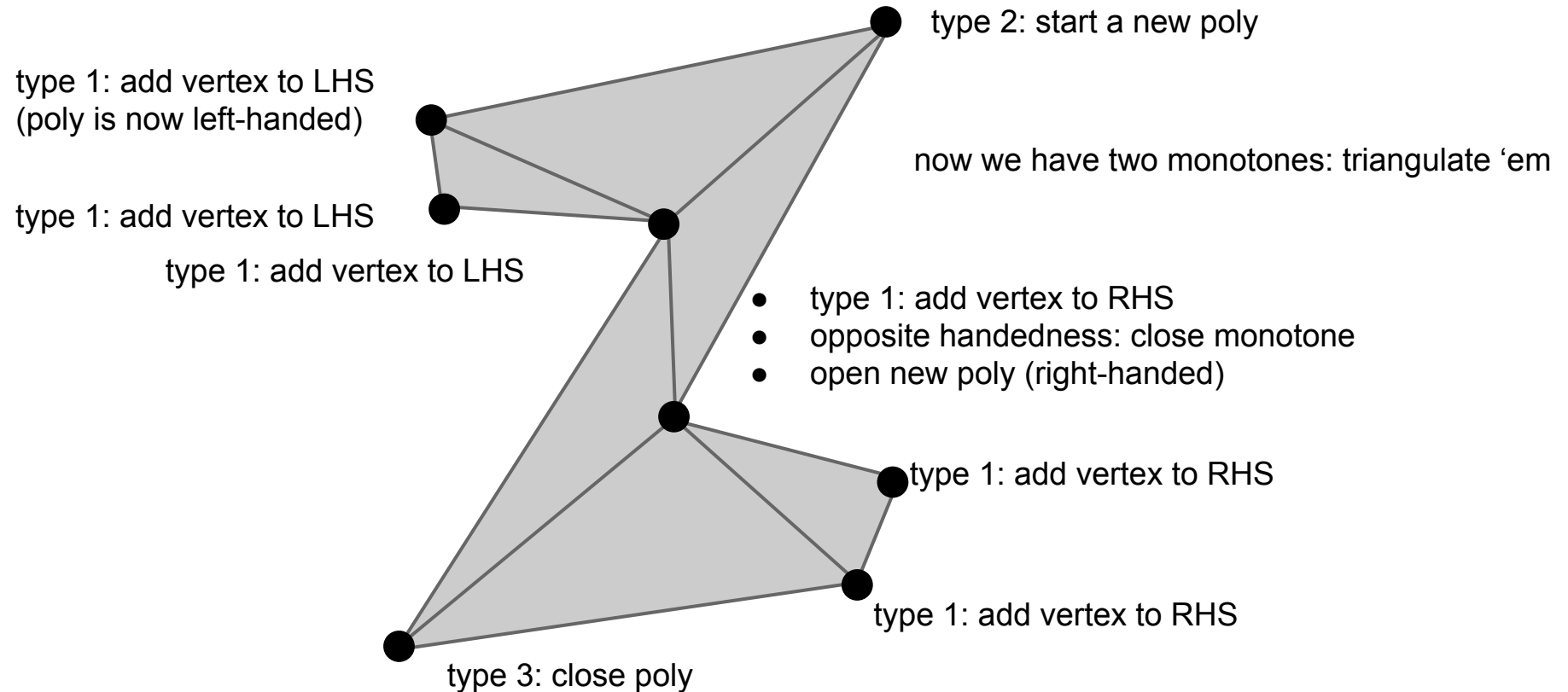


type 3

# Simple to monotone: da rules

- if incoming edge has left monotone
  - add current vertex to monotone's RHS
- if incoming edge has right monotone,
  - add current vertex to monotone's LHS
- if added on opposite side to handedness:
  - close monotone & start a new one
  - use previous & current vertex as the first two verts

# Simple to monotone: example



Fine print: left-handed monotones are triangulated bottom-up (simplifies code)



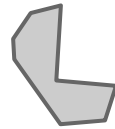
# Animals in the polygon zoo

- Complex

- self-intersections, holes

- Simple

- no self-intersections, no holes, concave



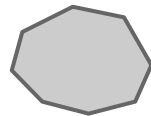
- Monotone

- monotonically incr; all pts one side, concave



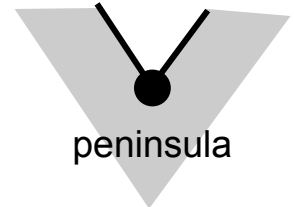
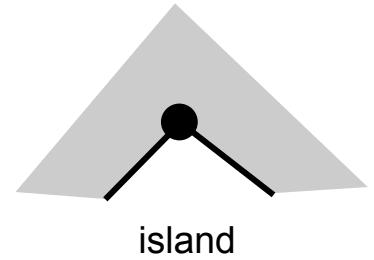
- Convex

- all angles  $\leq 180^\circ$



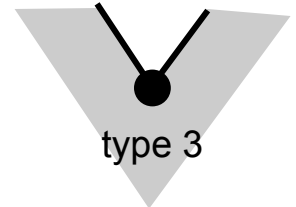
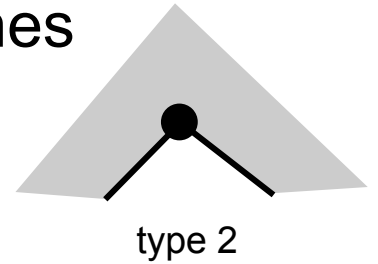
# What about islands / peninsulas?

- Island: a lone type-2 enclosed inside a larger poly (e.g., hole)
- Peninsula: type-3 that joins two previously-unrelated polys

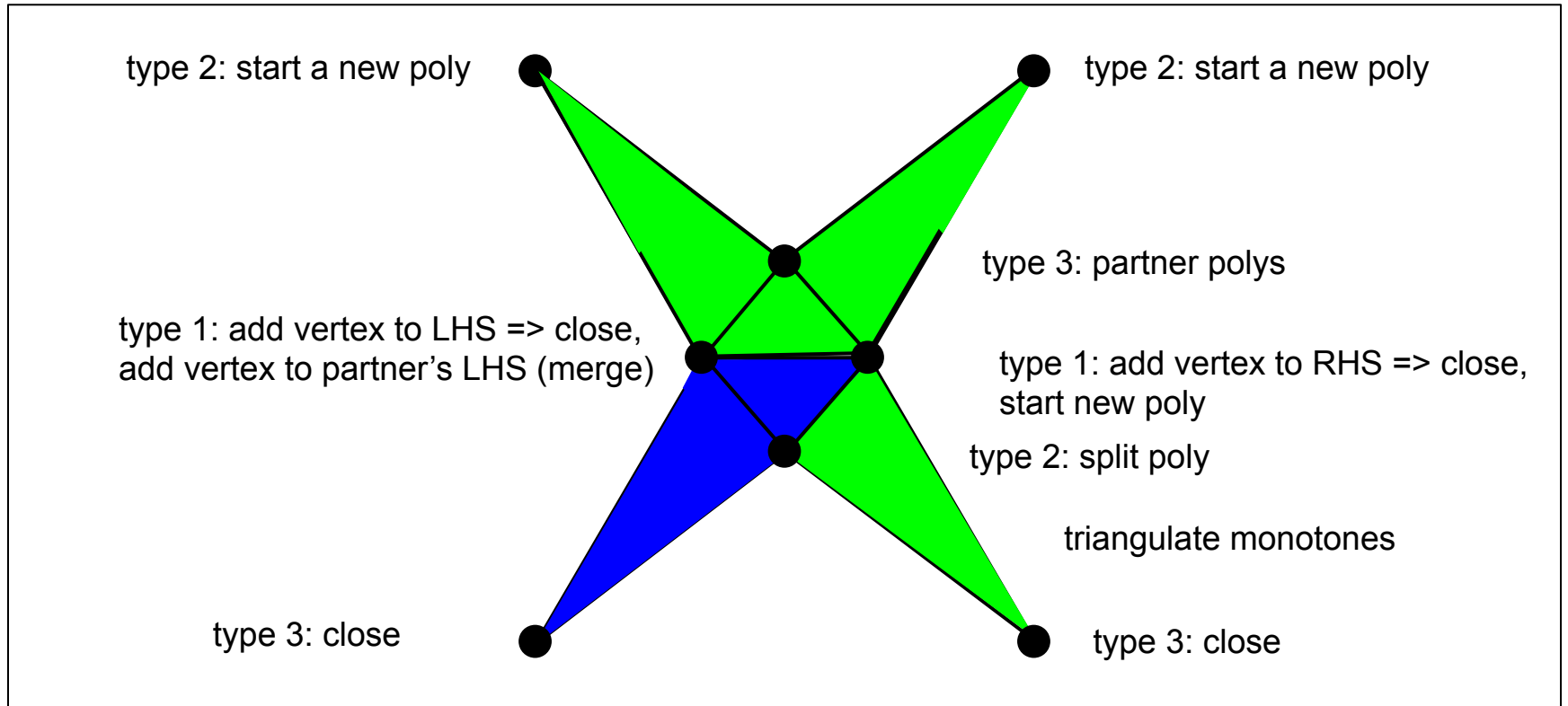


# Moar rules: splitting and merging

- Island:
  - if monotone above, split into two monotones
- Peninsula
  - merge left & right shapes via partnership
  - partnership ends at subsequent vertex
  - instead of starting a new monotone, add vertex to partner monotone



# Splitting and merging: example

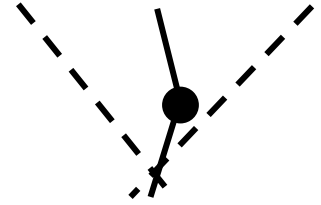


# Intersections

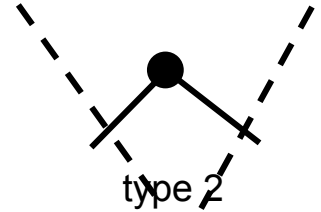
Seems  $O(N^2)$  in edges?

Bentley-Ottman (sweep line):

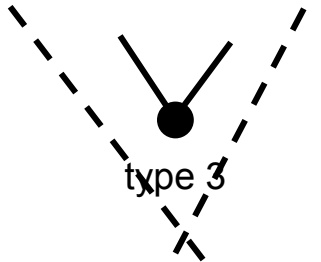
- key observation: all intersecting edges are neighbours in AEL
- check for intersections:
  - type 1 & 2: outgoing vs. enclosing edges
  - type 3: left enclosing vs right enclosing
  - $O((N + k) \lg N)$  for  $k$  intersections



type 1



type 2



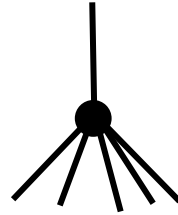
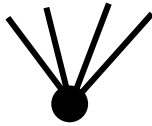
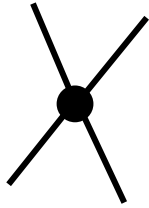
type 3

# Intersections, cont'd

- insert new vertex at intersection point
  - always below sweep line
- split edges
- check for coincident vertices; merge

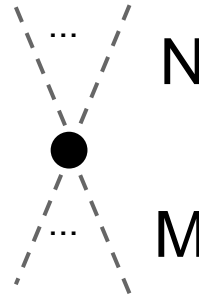
# Intersections

The vertex zoo just got bigger...



# Vertex zoo, generalized

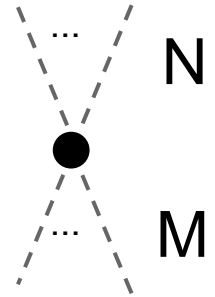
- N edges above, M edges below
- For AEL:
  - $N == 0 \Rightarrow$  type 2
  - $M == 0 \Rightarrow$  type 3
  - else  $\Rightarrow$  type 1





# Vertex zoo, generalized

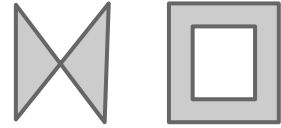
- When  $N > 2$ , close multiple polys above
- When  $M > 2$ , open multiple polys below
- poly membership is copied along leftmost & rightmost edges



# Animals in the polygon zoo

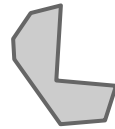
- **Complex**

- self-intersections, holes



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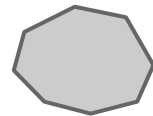
- **Monotone**

- monotonically incr; all pts one side, concave



- **Convex**

- all angles  $\leq 180^\circ$



# Pretty easy so far

- Nice computational geometry algorithms
- Fairly easy to implement
- $O(N \lg N)$

# Excrement, meet rotary ventilator

(fun with floating point)

- Bentley-Ottman requires exact math
- intersection is 5th order polynomial
- 24 bits mantissa => 160 bits :(

# Float sucks: what do we do?

Two schools:

- These algorithms don't work; give up
- Use arbitrary precision arithmetic

# Float sucks: what do we do?

Drop out of school:

- assume intersections are random;  
adjust mesh geometry to match
- inaccuracies are in ULPs, not pixels

# Excrement, meet rotary ventilator

Intersections can:

- change outgoing edge order within vertices
  - solution: intersection => check & reorder top vertex
- change edge order in active edge list
  - solution: intersection => check & reorder w/neighbours in AEL

# Excrement, meet rotary oscillator

Merging coincident vertices can:

- add edge above type 2  $\Rightarrow$  type 1
- remove last edge above type 1  $\Rightarrow$  type 2
- both can change enclosing edges and polys

Solution: on type change,

- redo find enclosing edges
- restart intersection checks



# Future work

- shipping in M44, Android, MSAA
- caching (threshold scale), shipping in M46
- enable in `<canvas>`
- alpha ramp AA (no MSAA required)
  - Desktop: Intel MSAA perf is terrible
  - 1-pixel-wide alpha ramp around geometry
  - batching

# References

- Fournier & Montuno, ["Triangulating Simple Polygons and Equivalent Problems"](#)
- Bentley, J. L.; Ottmann, T. A. (1979), ["Algorithms for reporting and counting geometric intersections"](#)
- Jean-Daniel Boissonnat and Franco P. Preparata. Robust plane sweep for intersecting segments. *SIAM Journal on Computing*, 29:1401-1421, 2000.

Questions?

# Performance results

<https://docs.google.com/spreadsheets/d/1S1hIb5hHCG04fP-9XeFgc73Zo07N8Hla5IYC07ju-Gw/edit#gid=0>

[https://docs.google.com/spreadsheets/d/1mMU39qbUt7LmQRUngqRJ0oL6sE7zwl8nElx3j\\_\\_bnek/edit#gid=0](https://docs.google.com/spreadsheets/d/1mMU39qbUt7LmQRUngqRJ0oL6sE7zwl8nElx3j__bnek/edit#gid=0)