# 5. Data Structures

2023-12-18 Section materials: jrsacher.github.io/cs50/

# Agenda

- The last week of C!
- Quick review of structs
- Linked lists
  - $\circ$  Nodes
  - $\circ$  -> operator
  - Comparison with arrays
- Stacks, queues
- Hash tables
- Trees
- Tries

### typedef and structs

```
typedef old-name new-name;
typedef char * string;
typedef struct car
  int year;
  char model[10];
  char plate[7];
  bool automatic;
car;
```

Accessing structs

// declaration
car mycar;

// field access
mycar.year = 2017;
strcpy(mycar.plate, "CS50");
mycar.automatic = true;

Accessing structs (with pointers!)

```
// declaration
car *mycar = malloc(sizeof(car));
// field access
mycar - year = 2017;
strcpy(mycar->plate, "CS50");
mycar->automatic = true;
// equivalent to (*mycar).automatic = false
```

# linked lists

### Arrays

Data stored sequentially in memory

- Good
  - Rapid, random access
  - Efficient storage of data
- Bad
  - $\circ$  Fixed size

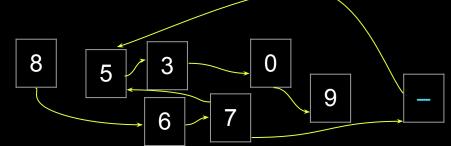
# "Growing" arrays



8	6 7	7 5	3	0	9	2
---	-----	-----	---	---	---	---



### Data stored separately, but connected



#### • Good

- $\circ$  Shrink and grow as needed
- Can insert or delete in arbitrary order
- Bad
  - No random access linear search
    - There are *technically* ways around this, but they get complicated.
  - Takes more memory to store the pointer takes extra space

# Linked Lists

- A linked list node is a special type of struct with (minimally) 2 properties:
  - Data of some type
  - A pointer to another node in the linked list.

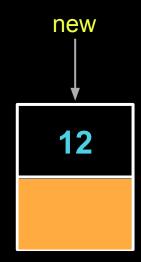
```
typedef struct node
{
    int value;
    struct node *next;
}
node;
```

### Operations needed to work with linked lists

- Creating a linked list when it doesn't exist.
- Inserting a new node into a linked list.
- Searching through a linked list to find an element.
- Deleting a single element from a linked list.
- Deleting an entire linked list.

# Making a new linked list

- Dynamically allocate space for a new node with malloc.
  - This returns a pointer to your newly created node.
- Make sure you didn't run out of memory.
  - Check that the pointer isn't NULL
- Initialize the value field.
- Initialize the next field (specifically, to NULL).



## In code ...

#### typedef struct node

// the value to store in this node
 int value;
 // the link to the next node in the list
 struct node \*next;

} node;

```
node *list = malloc(sizeof(node));
// Check that malloc succeeded
if (list == NULL)
{
    return 1;
}
// Add a value to the new node
list->value = 1;
list->next = NULL;
```

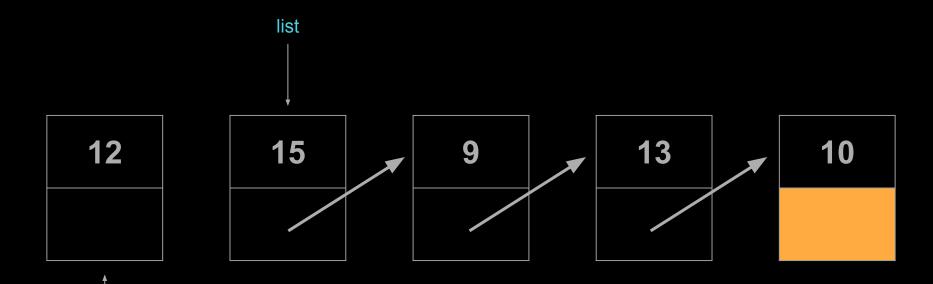
# Inserting

- Dynamically allocate space for a new linked list node.
  - Check to make sure we didn't run out of memory.
- Populate, insert node at the beginning of the linked list.
  - So which pointer do we move first? The pointer in the newly created node, or the pointer pointing to the *original* head of the linked list?
  - This choice matters!
- Return a pointer to the new head of the linked list.
- Note: can also insert into the middle or end of a linked list

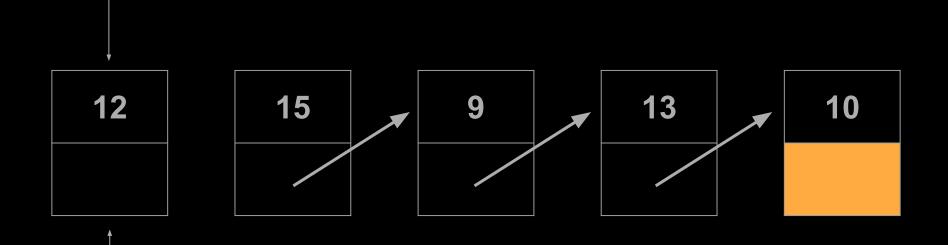
### In code ...

```
// Create a new node
node *n = malloc(sizeof(node));
// Check that malloc succeeded
if (n == NULL)
{
   free_linked_list(list);
   return 1;
}
// Add a value to the new node
n->value = 2;
// Link the new node to the list
n->next = list;
// move the list to point to the new node
list = n;
```

# list = n;

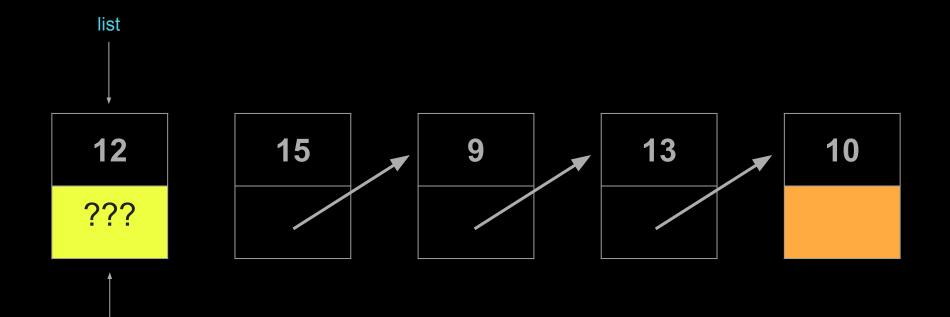


n



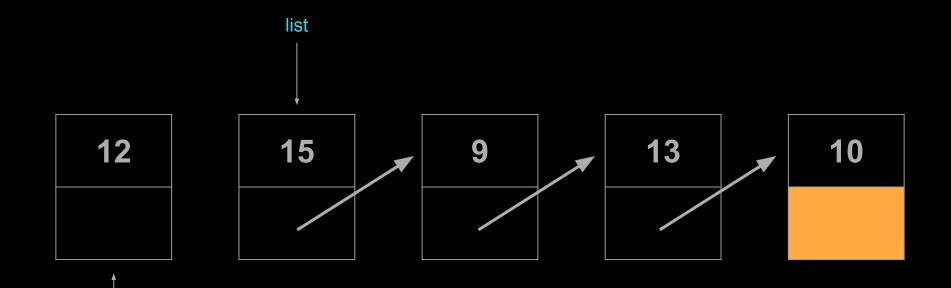
n

list



n

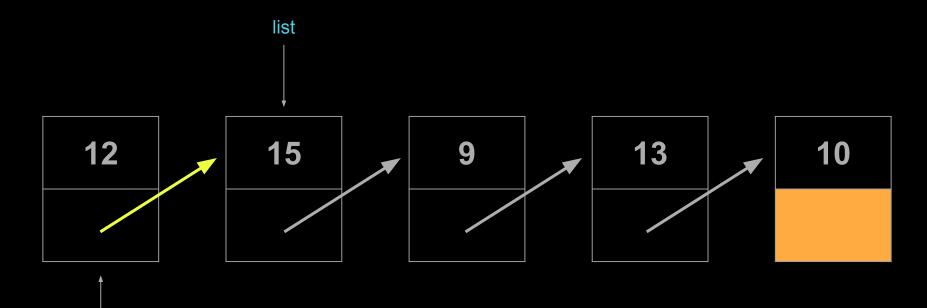




n

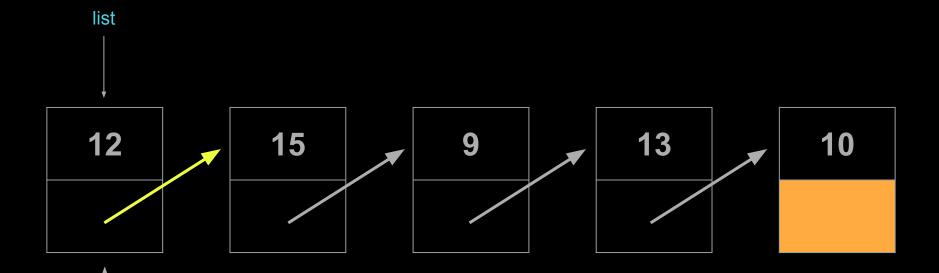
••

n->next = list;



n

list = n;



n



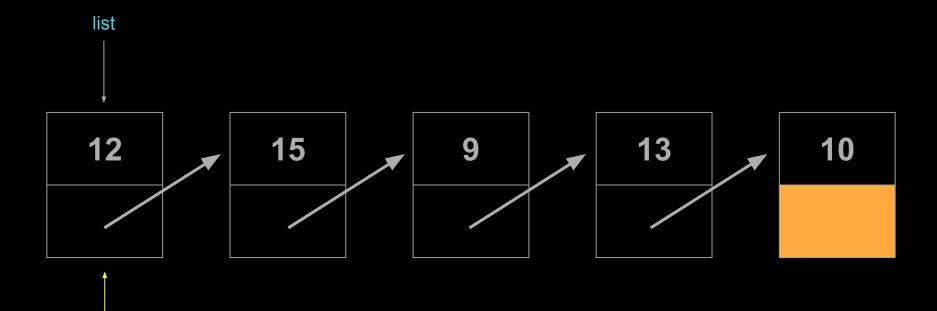
# Searching

- Make a traversal pointer to list's head (first element).
- If current node's value field is what we want, return true.
- If not, set the traversal pointer to the next pointer in the list and go back to the previous step.
- If you've reached the end of the list, return false.

### In code ...

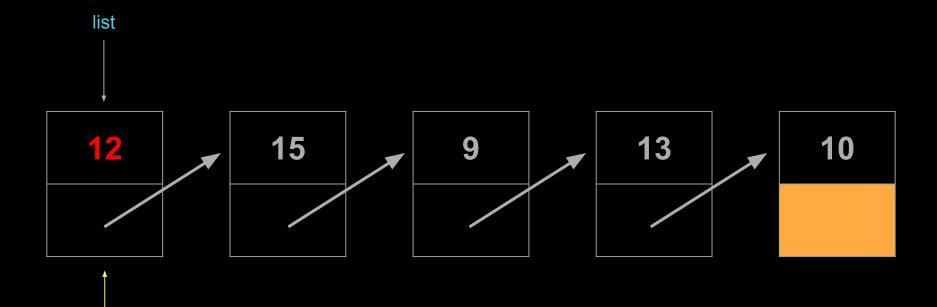
```
bool found = false;
node *tmp = list;
while (tmp != NULL && !found)
    if (tmp->value == n)
        printf("Found\n");
        found = true;
        break;
    tmp = tmp->next;
}
if (!found)
   printf("Not found\n");
}
```

```
node *tmp = list;
```



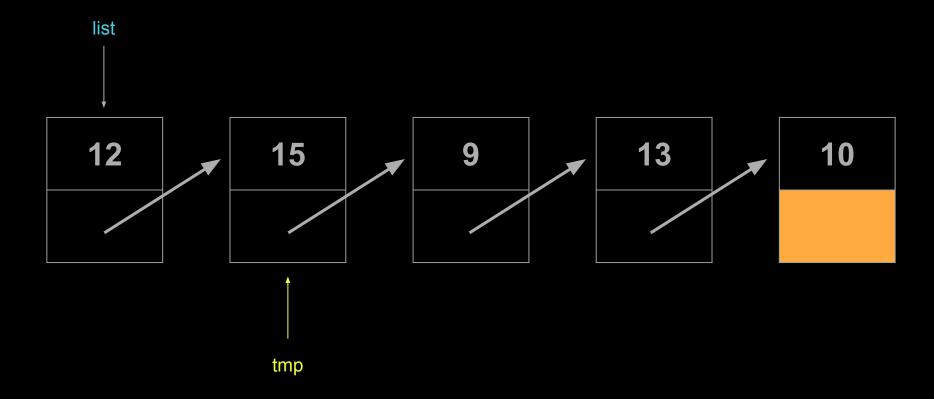
tmp

# tmp->value == 15;

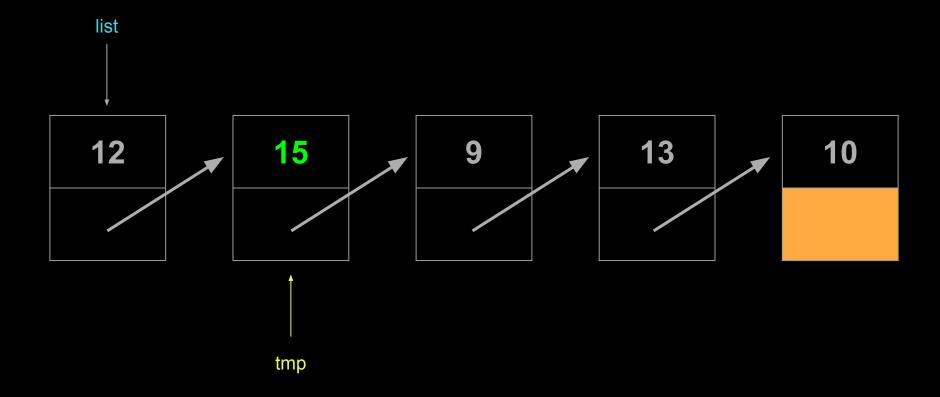


tmp

tmp = tmp->next;



# tmp->value == 15;



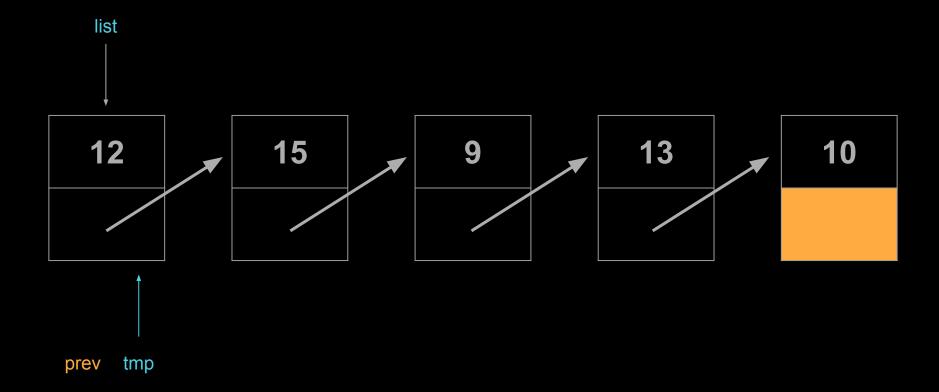
# Deleting

- A bit like a combination of searching and inserting
- If deleting from the middle, use 2 temp pointers to traverse linked list
  - Keep track of current and previous node
- When you find the value to delete
  - Point the previous node to the one pointed to by the current node to close the gap
  - $\circ$  free the current node

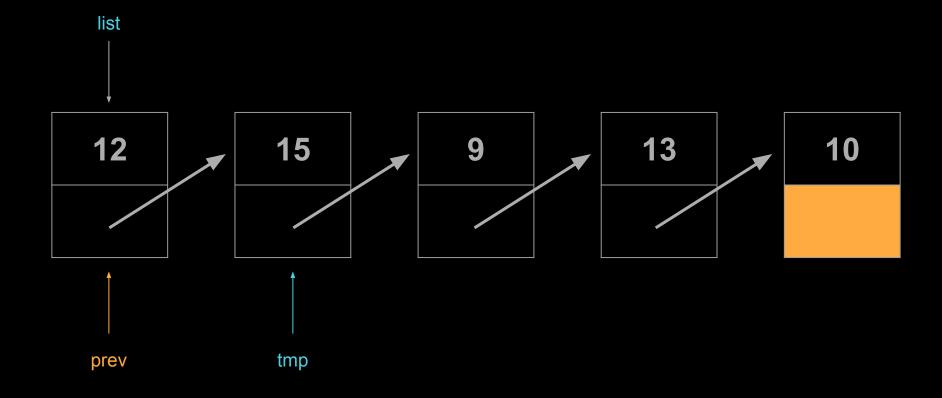
### In code ...

```
node *tmp = list;
node *prev = NULL;
while (tmp != NULL)
{
    if (tmp->value == n)
        if (prev == NULL)
            // we need to delete the first node
            list = tmp->next;
        }
        else
            // we need to delete some other node
            prev->next = tmp->next;
        free(tmp);
        break;
    prev = tmp;
    tmp = tmp->next;
```

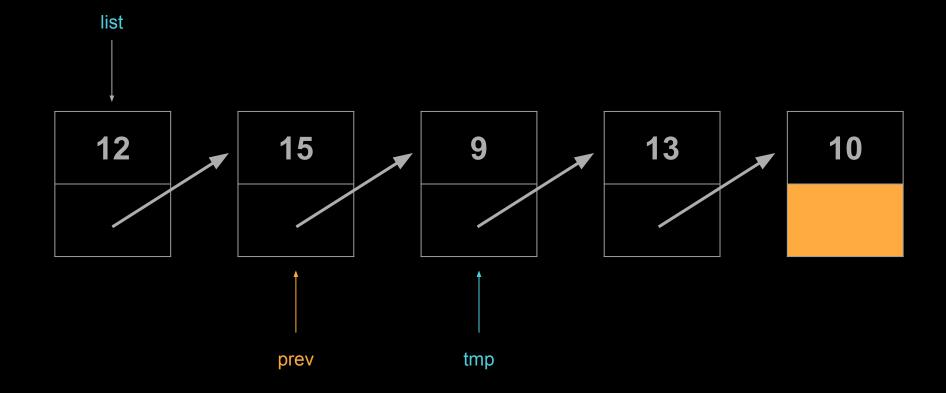
### node \*tmp = list; node \*prev = NULL;



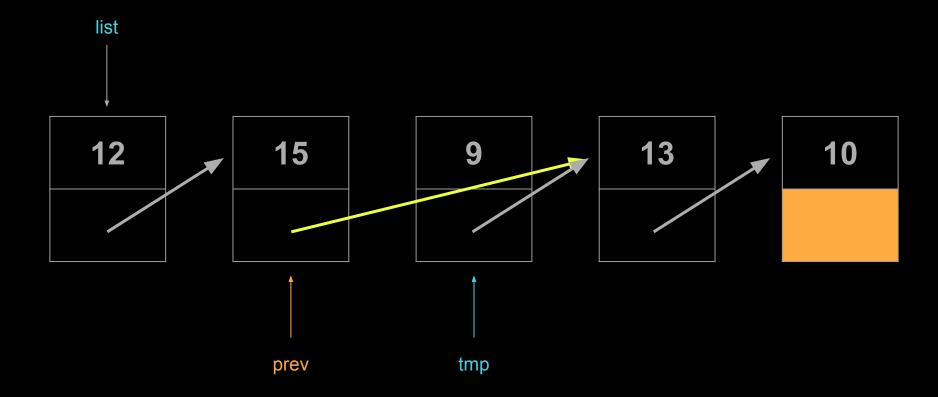
prev = tmp; tmp = tmp->next;



prev = tmp; tmp = tmp->next;



prev->next = tmp->next;

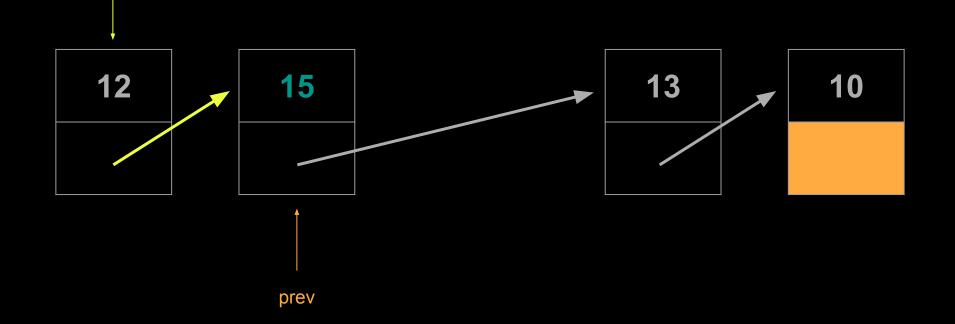


free(tmp);



### Note: -> operators can be chained! list->next->value

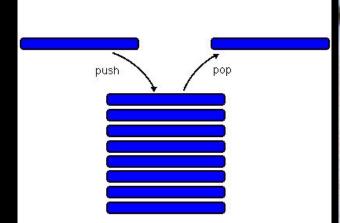
list



# stacks and queues

## Stacks

- Commonly an array or linked list
- Last-in-first-out (LIFO)
- 2 operations:
  - o Push
  - Pop





## Queues

- Also arrays or linked list, but have to keep track of "front" of the queue (previous and next pointers)
- First-in-first-out (FIFO)
- 2 Operations:
  - Enqueue
  - $\circ$  Dequeue



# Implementation as linked list nodes

```
typedef struct stack
{
    int value;
    struct stack *next;
}
stack;
```

(pretty much a regular linked list!)

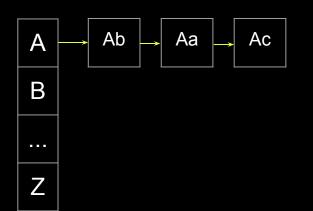
```
typedef struct node
   int value;
   struct node *next;
node;
typedef struct queue
    int count;
    node *front;
    node *rear;
queue;
```

Note: can be done in other ways, too!

## Hash tables

### An array of linked lists!

- Combines benefits of both
  - Random access to "bins"
  - Shorter linked lists to traverse
- Hash function determines which "bin" to use
  - Must be deterministic, only return a fixed number of choices



## Hash functions

- A hash function finds the appropriate "bucket" for data
- A good hash function should:
  - Use only the data being hashed
  - Use all of the data being hashed
  - Generate the same result every time given same input; no randomness!
  - Uniformly distribute data
  - Generate very different hash codes for very similar data.

### ...ahem

```
const int N = 1000000
// hash function described in http://www.cse.yorku.ca/~oz/hash.html
unsigned int hash(const char *str)
ł
   unsigned int hash = 5381;
   char c;
   while ((c = *str++))
   {
       hash = ((hash << 5) + hash) + tolower(c);
   }
   return hash % N;
}
```

Note: use a **big** number for N!

#### hash.c

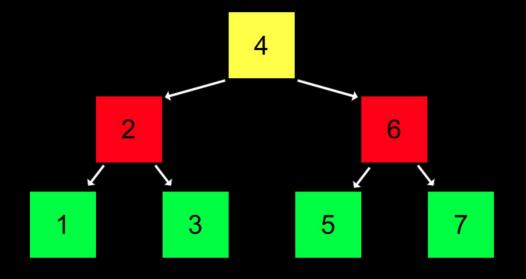
## Collisions

- Unless your hash function is perfect, collisions will happen
- Solution 1: Linear probing (not ideal in our case)
  - If we have a collision, try to place the data in the next consecutive element of the array instead
  - Try each consecutive element until we find a vacancy.
  - That way, if we don't find it right where we expect it, it's hopefully nearby.
- Solution 2: Chaining (with linked lists)
  - Each element of the array is now a pointer to the head of a linked list.
  - For a collision, make a new node and add it to the chain at that location.
  - $\circ$   $\,$  The item is where it's expected, just in a list

# trees

## Binary search tree

- Similar to a linked list, but instead of **next**, it has 2 pointers: **left** and **right**
- To be effective:
  - Data must be sorted
  - The tree must be balanced (~same number of entries on each side)
- Gives the speed of binary search with the flexibility of a linked list



# tries

## Tries

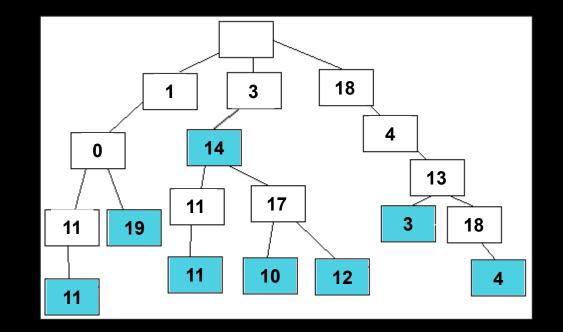
- Tries have unique keys, with values as simple as a bool to tell you that the data exists in the structure
- No collisions no two pieces of data have the same path
- Insertion, deletion, and lookup all O(1)!
  - In reality, takes up to *n* steps, where *n* is the height of the trie
- Tradeoff is space!

```
Trie struct for dictionary use
```

```
typedef struct node
{
    bool is_word; // 1 byte
    struct node *children[27] // 8 bytes * 27 = 216
}
node;
```

- 27 needed for 26 letters + ' (apostrophe)
- Word is never "stored" as a whole -- only the path

# Graphical trie



# Questions?