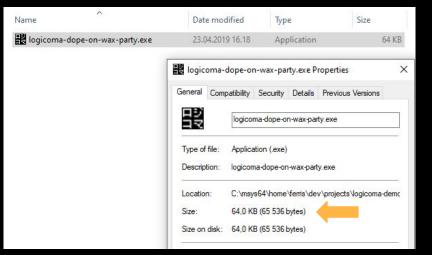
modern(ish) 64k intro compression jake "ferris" taylor / logicoma

• demo that fits in 65536 bytes (often less because lazy)



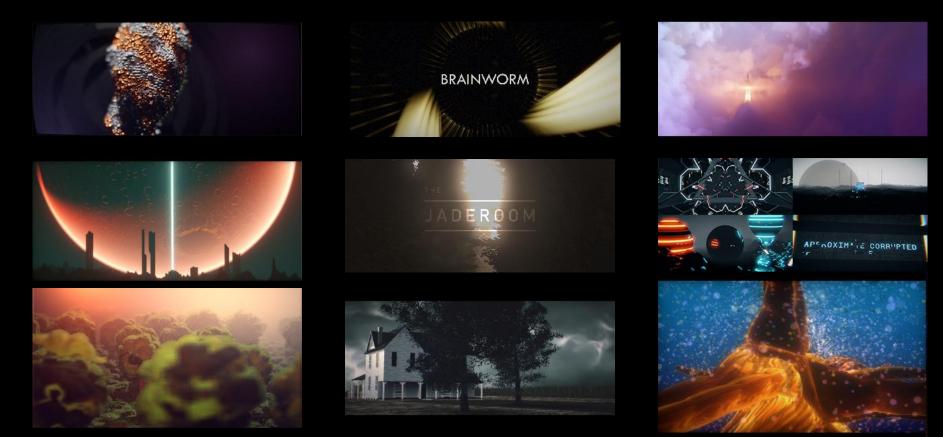
- demo that fits in 65536 bytes (often less because lazy)
- single executable, no external media (except OS/drivers)

Name	Date r	nodified	Тур	e	Size	
🔡 logicoma-dope-on-wax-party.exe	23.04.2019 16.18		Application		n 64 I	
	B logicon	na- <mark>dope-o</mark> i	n-wax-pa	rty.exe P	roperties	×
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			coma-dope-on-wax-party.exe			
	Type of file	e: Applica	ation (.exe)			_
	Description	n: logicon	na-dope-or	n-wax-par	ty.exe	_
	Location:	C:\msy	s64\home	\ferris\de	v\projects\logicoma-d	emc
	Size:	64,0 K	B (65 536 I	bytes)		
	Size on dis	sk: 64,0 K	B (65 536 I	bytes)	•	

- demo that fits in 65536 bytes (often less because lazy)
- single executable, no external media (except OS/drivers)
- coolest demoscene category imo (totally not biased)

Name	Date mod	dified	Type	Size
🔡 logicoma-dope-on-wax-party.exe	23.04.201	9 16.18	Application	64 KB
	coma-	01		×
	Com	pa	0	ns
	비와 기자	logicom		
	Type of file:	Applicatio	n (.exe)	
	Description:	logicoma-	dope-on-wax-party.exe	•
	Location:	C:\msys6	4\home\ferris\dev\pro	jectsVogicoma-demc
	Size:	64,0 KB (65 536 bytes)	
	Size on disk:	64,0 KB (65 536 bytes)	

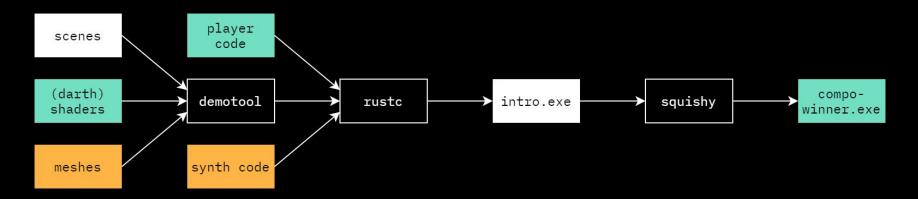
here's what some of them look like



here's what some of them look like

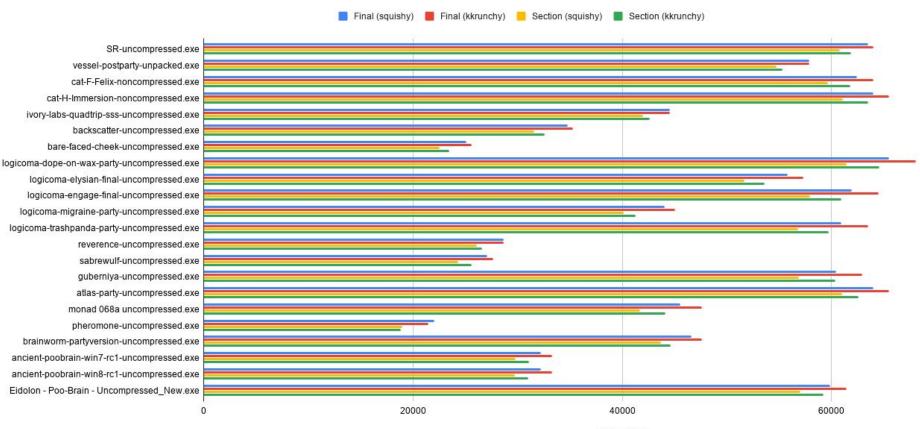


- our (logicoma's) executable compressor
- developed since 2016
- specifically built for 64k
 - much heavier compression engine than 1k/4k/8k
- http://logicoma.io/squishy/



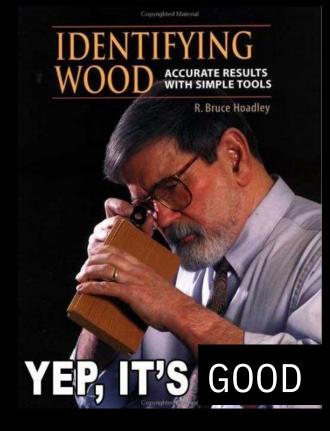
• is it good?

squishy vs kkrunchy metrics

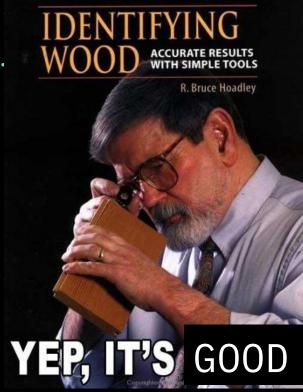


• is it good?

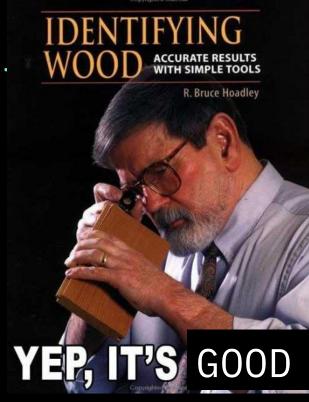
• is it good?



- is it good?
- could be better..



- is it good?
- could be better..
- still wip :)



full disclosure..

- squishy was most focused on improving the compression engine
- all other techniques are icing on the cake
- as long as the main engine kicks butt!

full disclosure.

- to me, this is the boring part • and often covered in other talks!
- but overview is necessary
- so I'll describe some stuff

executable compression
(for reals this time)

- .exe -> .exe (instead of .whatever -> .ziphead)
- OS needs to be able to execute it!
- most metadata not needed
- some things can be folded/overlapped
- function imports?
- resources need to be uncompressed (but not all!)

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DOS Header

(0x3C) Pointer to PE Head

ignature 0x5A4D

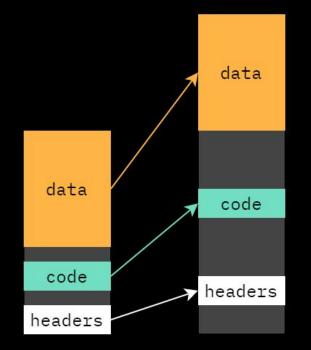
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https://en.wikipedia.org/wiki/Portable_Executable#/media/File:Portable_Executable_32_bit_Structure_in_SVG_fixed.svg

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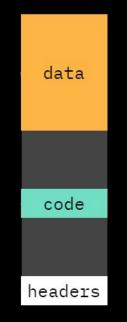
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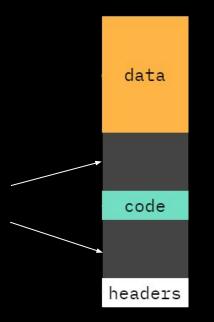
data code headers

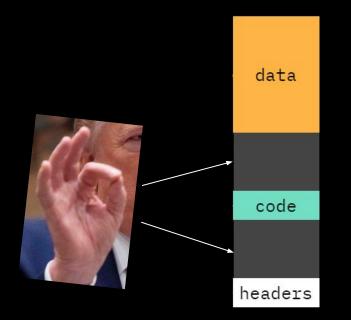


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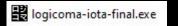
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- resolve at runtime using above fn's

• some special handling for resources

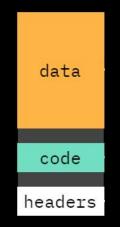


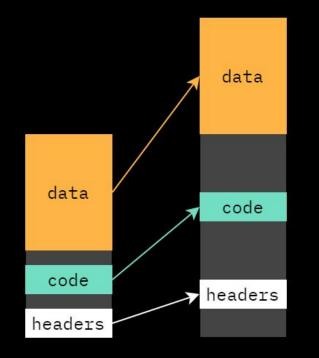


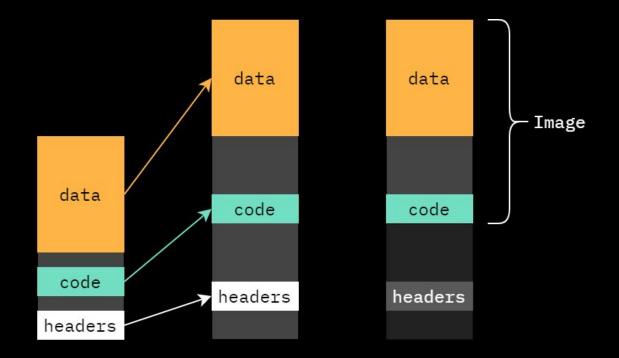


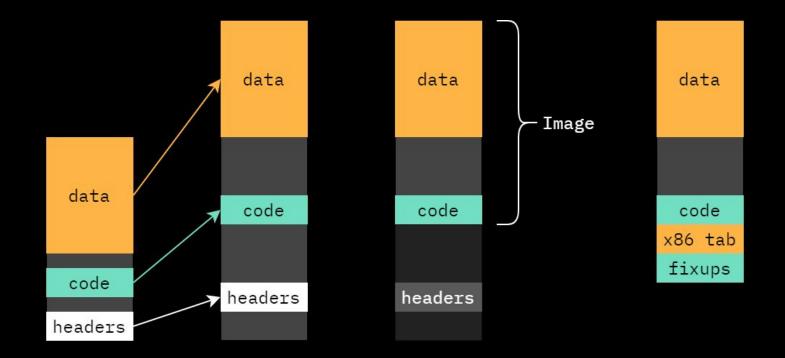


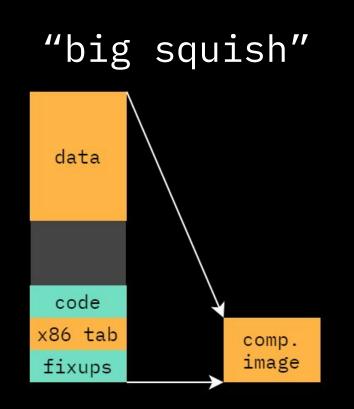
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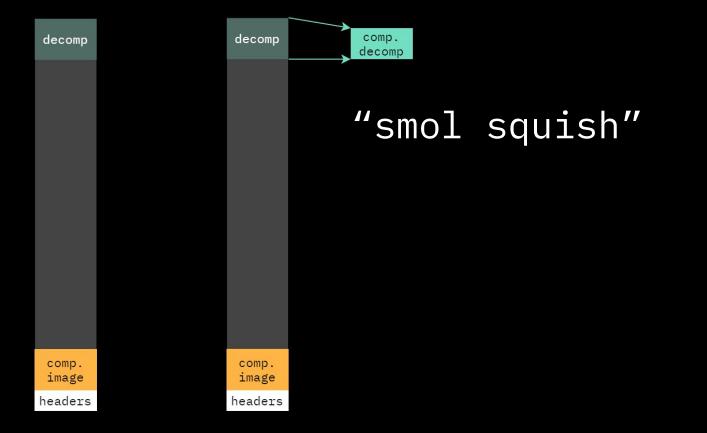


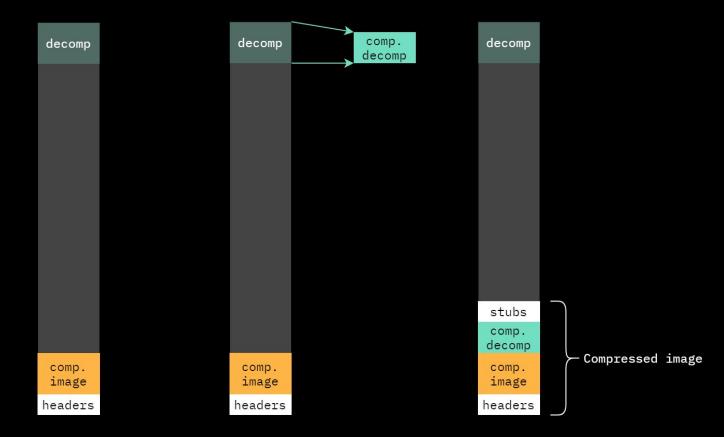


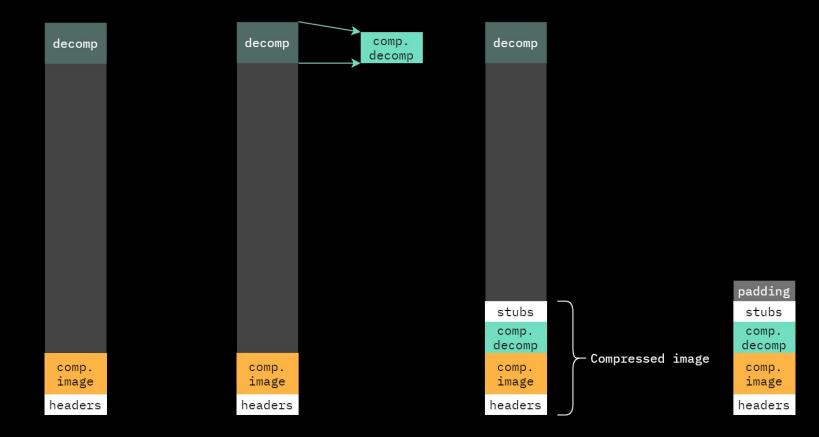


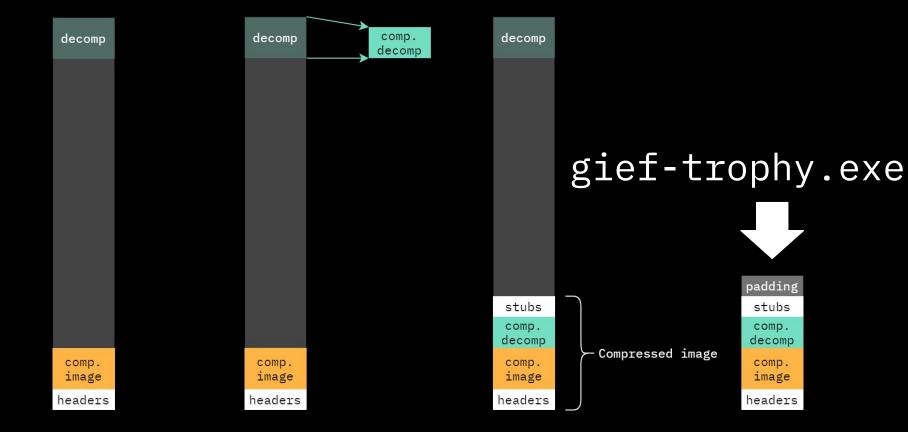




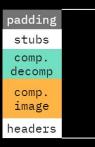












padding stubs comp. decomp comp.

comp. image

padding

stubs

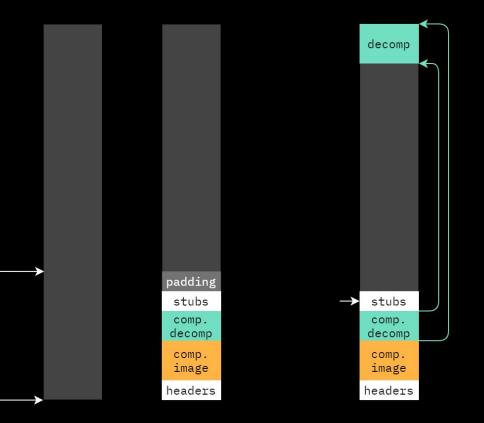
comp.

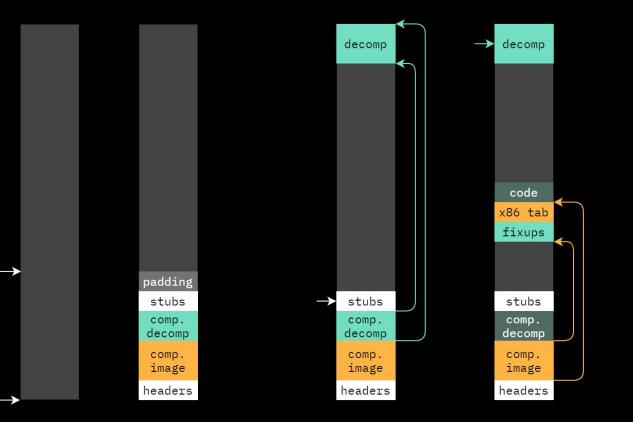
decomp

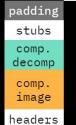
comp.

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padding stubs

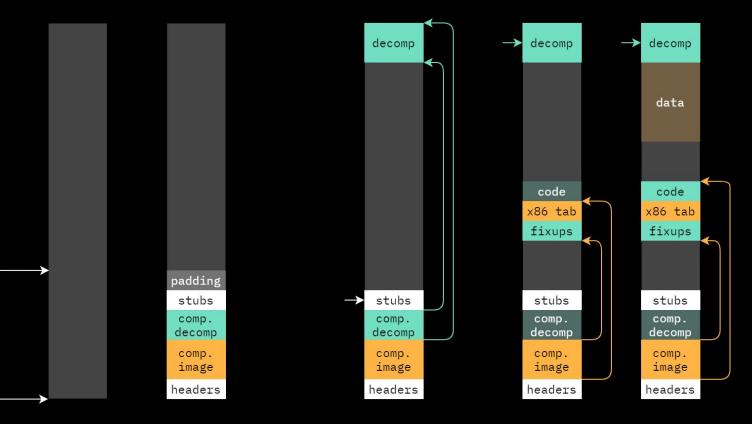
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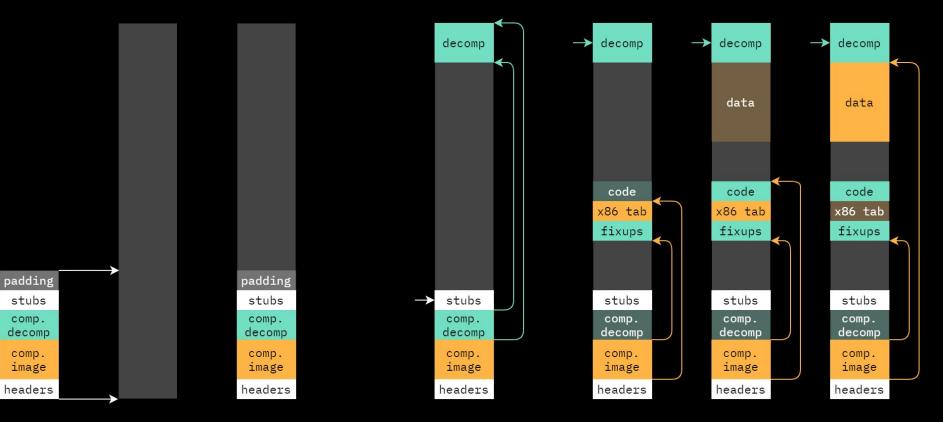


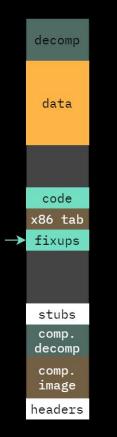
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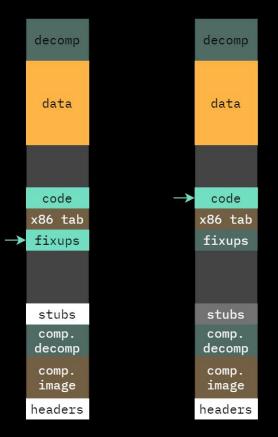
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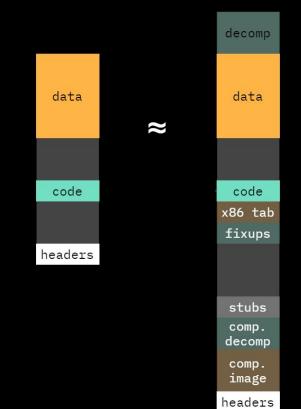
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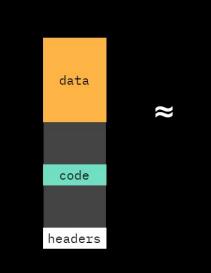
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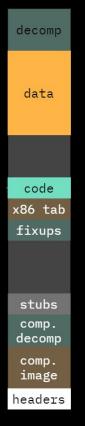














executable compression (final notes)

- we don't do very advanced slicing/header packing
- no hash import, minimal overlapping
 - largely compatibility measures
 - we can afford this in 64k
 - compatible with future Windows loaders = better user experience
- compression engine pulls most of the weight
- maybe I'll cave and add some hackier stuff optionally anyways eventually :)

compression 101

quick disclaimer

- two compressors in squishy
 - \circ both have statistical components
- only have time to talk about the big one
 - \circ $\mbox{ more "pure" in the theoretical sense anyways }$
 - \circ everything in this section still applies to both

compression 101

- start with a symbol alphabet
 - \circ could be any length >= 1
 - o eg. {0, 1} for bits, {0, 1, .. 255} for bytes, {A, B, C}, etc.

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• find a new representation which requires fewer bits but conveys the same information



most real-world data contains statistical redundancy
 put simply: some symbols are more common than others

• we can exploit this by making a statistical model of the string we want to compress

- let's take an example, our string from earlier: AABC
- now we'll determine the frequency of each symbol
 this just means count them!

- how many A's are there in our string?
 o 2, so f(A) = 2
- how many B's are there in our string?
 1, so f(B) = 1
- how many C's are there in our string?
 - \circ 1, so f(C) = 1

- let's take an example, our string from earlier: AABC
- now we'll determine the probability of each symbol
 this just means divide the frequencies by the total string length!

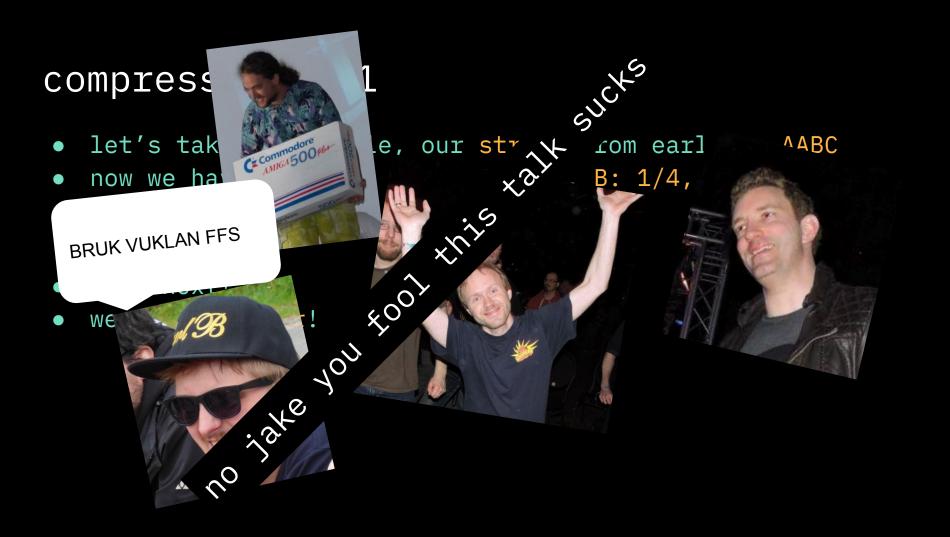
- len = 4 (chars)
- p(A) = f(A) / len = 1/2
- p(B) = f(B) / len = 1/4
- p(C) = f(C) / len = 1/4

note how these are normalized, i.e. they sum to 1

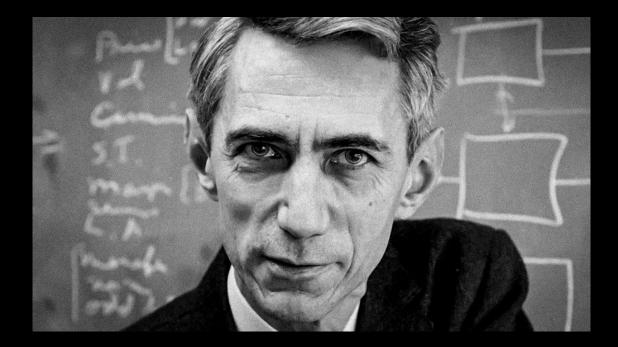
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• what next?

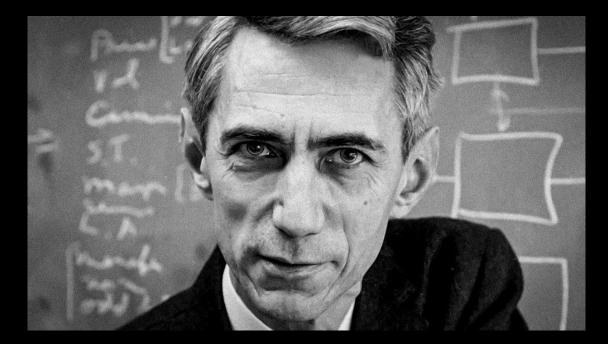


• meet claude shannon



- meet claude shannon
- look at that hot piece of man(tropy)

 ok this pic is haunting af but hear me out



- claudy with a shans of meatballs over here did this really cool thing
 - he actually did a butt ton of awesome stuff!!!

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 where p is the probability of the symbol, as discussed earlier

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- the optimal code length for a symbol is -log2(p) bits
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 - for A: $-\log^2(1/2) = 1$ bit
 - for B: $-\log^2(1/4) = 2$ bits
 - for C: $-\log^2(1/4) = 2$ bits

- the optimal code length for a symbol is -log2(p) bits
- our earlier model model: p = { A: 1/2, B: 1/4, C: 1/4 }
 - for A: $-\log^2(1/2) = 1$ bit
 - for B: $-\log^2(1/4) = 2$ bits
 - for C: $-\log^2(1/4) = 2$ bits

- an [entropy] coder codes symbols using probabilities, such that each symbol is represented using the optimal number of bits as calculated above
 - \circ ^ THAT's the kind of coder we need!

- back to our example string from earlier: AABC
- and our model: p = { A: 1/2, B: 1/4, C: 1/4 }
- and a coder **c** with the following interface:

```
class coder:
    output_string: string
    coder():
        output_string = empty;
    method encode(symbol, model):
        /* ignore impl for now */
```

- back to our example string from earlier: AABC
- and our model: p = { A: 1/2, B: 1/4, C: 1/4 }
- and a coder c
- we simply invoke our coder for each symbol, and it handles the rest! (assuming our model is correct)

```
fn encode(string, model):
    c = new coder();
    for s in string:
        c.encode(s, model);
    return c.output_string;
```

- back to our example string from earlier: AABC
- and our model: p = { A: 1/2, B: 1/4, C: 1/4 }
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string: AABC

output_string: empty

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output_string: 0 0 10 11 = 001011 (6 bits!)

fn encode(string, model):
 c = new coder();
 for s in string:
 c.encode(s, model);
 return c.output_string;

fn decode(output_string, original_string_len, model):
 original_string = empty;
 d = new decoder(output_string);
 for i in 0..original_string_len:
 original_string.append(d.decode(model));
 return original_string;

- contrived example, but it illustrates key things:
- general statistical coding algorithm for compression
 - model determines per-symbol probabilities
 - coder faithfully encodes symbols with # of bits determined by the model probabilities
- model/coder are completely decoupled!

- many kinds of coders
- we saw an example of a prefix code "implementation"
 - huffman coding belong to this family
 - limited to fixed bit widths due to direct symbol replacement
 - \circ real-world probabilities are rarely powers of 2!
 - eg. the alphabet { A, B, C } with string ABC gives us 1/3 prob for each symbol, and a total length of 4.75 bits (approx.)

- there exist coders that output fractional bits
- you won't BELIEVE this ONE SIMPLE TRICK!
 - coder keeps some internal state representing fractional bits
 - \circ consider a symbol that should be represented with 1/10 bits
 - for every 100 of these symbols, the coder keeps track of fractional bits and only outputs 10 bits
 - o it's just averages!
 - \circ hidden behind (de)coder per-symbol interface

- I wish I had more time to talk :(
 - arithmetic/range coders
 - asymmetric numeral systems family

- tl;dr: coding is a well-understood, largely solved problem
- squishy uses a simple binary range coder
 - good fractional precision, fast (for binary models)
 - **rABS** was also experimented with, no compelling advantages

• modeling, however, is the hard part!

- getting a model right is really hard
 - \circ often requires intimate knowledge of the data
- 64k's can contain all kinds of different data
- we need a good general-purpose model
 - \circ needs to basically handle anything
 - perhaps tuned to shaders/text a bit these days
- one thing in common: x86 code
 - specialized modeling for this in squishy
 - I tried not to do this for a _long_ time, eventually caved!
 - o it's that important!

• we saw a ternary, static, order 0 model previously

• we saw a ternary, static, order 0 model previously ^ 3 symbols (A, B, C)

• we saw a ternary, static, order 0 model previously ^ same model used to code whole string

• we saw a ternary, static, order 0 model previouslywe'll get to this part :) ^

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for speed and simplicity, as we'll see

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- our model should adapt to changing statistics throughout the data for better compression

 \circ this also means we don't have to store a model in the compressed file

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- different models are good for different patterns in the data
- so let's run several models in parallel and mix results somehow

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 - PAQ
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- luckily we only compress a few hundred kb's so whatevs

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 - hence "order 0"
- consider an english sentence:
 - \circ $\,$ the cat kicked the dog in the face
- an order 0 model doesn't care about placement of words
 - \circ might as well have been cat dog face kicked in the the the
 - \circ same per-word (symbol) count for the whole string as above

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- but we know that context matters!
- a better model would give nouns (cat, dog, face) higher probabilities after articles (the)
 - \circ and loads of other rules like this
- a context that includes one symbol before the current symbol is an order 1 context
 - just like markov chains
 - \circ can have order 2, 3, ... 100 if you want
 - contexts become more sparse as order increases
 - \circ they also get YUGE!

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 - o eg. model_state = states[context_hash]
 - fetch prediction from model_state, update after symbol is seen

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 - different eviction strategies possible (eg. LRU)
 - squishy uses a 32MB 4-way cache table with LRU
 - optimized for cache line alignment and compression efficiency

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- this is fine since models can be arbitrary history functions
- note that many useful contexts are sparse!
 - eg. order 2 context with the byte that was 4 bytes ago and the one that was 8 bytes ago, instead of the 2 previous

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- so yeah, our job is to build models that predict bits, and use them to model our data in different contexts
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 - eg. only represent p1, since p0 = 1 p1
- scale by const power of 2, eg. 4096
 - 0 (0, 4096) instead of (0, 1)
 - this is what's used in squishy/PAQ/kkrunchy
 - for final probs at least, most models have more bits
 Internally; I want to increase this at some point!

• baby's first model:

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fn prob():
 return 1337;

• stationary context model

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^ current prediction, init to 1/2 (usually)

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^ # of bits seen so far, init to 0

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- two variables: prob, bit_count
- update rule biases prob in inverse proportion to bit_count

• stationary context model

```
fn init():
    prob = 2048;
    num_bits = 0;
```

fn prob():
 return prob;

```
fn update(bit):
    prob += (bit * 4096 - prob) / (num_bits + delta);
    if num_bits < limit: num_bits++;</pre>
```

- stationary context model
- two variables: prob, bit_count
- update rule biases prob in inverse proportion to bit_count
- delta and limit are tunable parameters
 - \circ best values depend on data
 - static, hand-tuned in squishy based on test corpus performance
 - maybe exposed in future versions

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- initially learns quickly, then becomes static (hence "stationary")
 - eventually becomes adaptive again as bit_count saturates, but learns slowly at that point

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- what kind of data might it model?

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- update **both** table entry and model state after symbol
 - \circ order matters here!
- possibly interpolate/update adjacent entries

• indirect context model

```
fn init():
    counts = 0 | 0;
    last_n_bits = 0;
```

```
fn prob():
    return prediction_table[counts | last_n_bits];
```

```
fn update(bit):
    prediction_table[counts | last_n_bits] += ...;
    counts += ...;
    last_n_bits = (last_n_bits << 1) | bit;</pre>
```

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 - ^ consider the end of each 3-symbol block

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^ assume an order-2 context, so the last two symbols are context

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 - \circ 000 001 000 001
 - ^ "after we see 00, there's a repeating 0, 1, 0, 1 pattern"

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- number of history bits to track, prediction table size, and update parameters are all tunable
 - \circ again, static in squishy, tuned on corpus
 - maybe exposed someday

• many more primitive model types available!

- many more primitive model types available!
 - run models for long strings of the same symbol
 - match models for higher-order contexts
 - variable-order models (eg. PPM, DMC, CTW)
 - \circ this is a fun place to be creative!

wow is this guy really still talking

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- short answer: however we want!
- let's look at some options!

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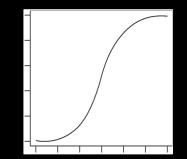
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- a is tunable
- not particularly useful
 - basically saying one model is always better than the other
- but we can do loads of stuff here!

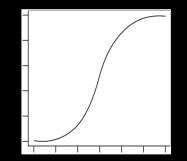
- linear mixing with fixed weights
- exactly what it sounds like
 - \circ p(X) * a + p(Y) * (1 a), a in [0, 1]
- a is tunable
- not particularly useful
 - \circ basically saying one model is always better than the other
- but we can do loads of stuff here!
 - average?
 - weighted average?
 - be creative!

• logistic mixing

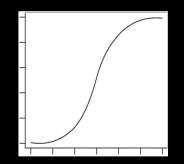
- logistic mixing
- remap input predictions on a logistic curve



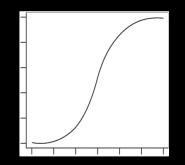
- logistic mixing
- remap input predictions on a logistic curve
 - \circ more precision at ends
 - \circ better scaling for logarithmic source coding equation



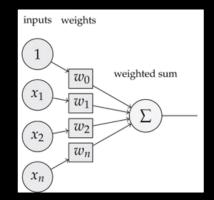
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- weighted sum remapped predictions

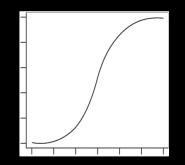


- logistic mixing
- remap input predictions on a logistic curve
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- weighted sum remapped predictions
- remap sum back to linear domain

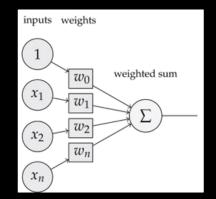


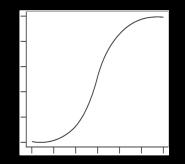
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- this is a perceptron



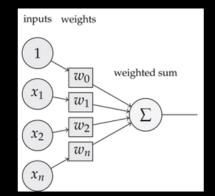


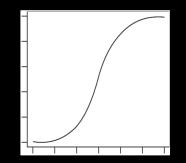
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- performs very well!





- logistic mixing
- remap input predictions on a logistic curve
 - \circ more precision at ends
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- weighted sum remapped predictions
- remap sum back to linear domain
- this is a perceptron
- performs very well!
- can select different weights with a context
 - squishy uses several!





wow yeah he's seriously still talking

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 - SSE/APM
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- we have loads of freedom to mix/match/adjust predictions
- how about postprocessing?
 - SSE/APM
 - ISSE
- multiple mixer stages?
 - you betcha!
- enabling multiple models for different data?
 - o sure, why not?
 - \circ this is part of how we model x86 in squishy actually

- based on PAQ7 with hand-tuned models
 - increased precision everywhere
 - several improvement ideas from PAQ8/ZPAQ/kkrunchy/etc
 - some special x86 stuff (later)
- ~2017/2018 focused on a ZPAQ-like VM
 - genetic algorithm used to "grow" architectures
 - never reached suitable performance :(
 - I didn't quite understand multiple mixer weight contexts at the time
 - models weren't as good as they are now
 - \circ will probably try this again sometime!
- not set-in-stone
 - I want to experiment more :)

• context models

- single 32MB cache table
- each entry (16 bytes) contains:
 - cache tag (4 bytes)
 - a stationary/direct model (4 bytes)
 - an indirect model (4 bytes)
 - a run model (4 bytes)
- 4-way associative cache
 - each bucket is 16 bytes * 4 = 64 bytes
 - same size as cache lines, aligned mem alloc

• context models

- o 21 contexts used (in data sections)
- hand-picked, static descriptions
- combined previous bytes (and bits of those bytes) selected by masks
- each context modeled by one cache table entry (3 models/predictions)
- o 21 * 3 = 63 context model predictions (in data sections)
- single const model
 - \circ also hand-tuned
 - o 63 + 1 = 64 predictions (in data sections)
- 8 match models
 - each with increasing context orders
 - o 64 + 8 = 72 total predictions (in data sections)

- logistic mixing
- first stage
 - all 72 predictions are mixed 8 different times
 - \circ $\,$ each time with weights selected from different contexts $\,$
 - 1 static context (order 0), same weights every time
 - 5 byte history contexts with increasing order (orders 1-4 and 8)
 - 1 bit history context
 - 1 weird, custom context (some match model state and other stuff)
 - o key here is to mix/match stuff!
 - pulls a LOT of weight!
- second stage
 - 8 mixed outputs mixed again by second stage with static context

- logistic mixing
- both stages mixed with 16 bit * 8 lane SIMD
- not super fancy, mostly SSE2 with SSSE3 horizontal sums
- basically the only SIMD in the whole thing

• APM stages

- final mixer output adjusted by 3 APM stages in serial
- each with increasing context orders
- static, linear weights
- not a huge difference after heavy modeling, but pays for itself

- final output clamped to [1, 4095] and sent to coder
- lots more possibilities
- this is what worked so far
- future squishy versions will likely do different stuff
 - \circ or not, who knows
 - \circ we have enough tooling/demos to make already as it is!

• tried to avoid, couldn't to be competitive!

- tried to avoid, couldn't to be competitive!
- e8/e9 filter
 - this will be replaced shortly due to false positives
 - all experiments with fancier cache schemes so far help code compression, hurt in total, needs further work
- need something more comprehensive

- main idea: leave code in-place
 - don't reorder like kkrunchy in case there are useful correlations
- use same models as in data sections
- on-the-fly state machine disasm

00602161	Ofb6460e	movzx eax, byte ptr [esi + 0xe]
00602165	db45f8	fild dword ptr [ebp - 8]
00602168	dd05b0997800	fld qword ptr [0x7899b0]
0060216e	8945f8	mov dword ptr [ebp - 8], eax
00602171	8b450c	mov eax, dword ptr [ebp + 0xc]
00602174	dcf9	fdiv st(1), st(0)
00602176	0300	add eax, eax
00602178	33db	xor ebx, ebx
0060217a	8945£0	mov dword ptr [ebp - 0x10], eax
0060217d	d9c9	fxch st(1)
0060217f	d95dfc	fstp dword ptr [ebp - 4]
00602182	db45f8	fild dword ptr [ebp - 8]
00602185	de£1	fdivrp st(1)
00602187	dc0d189a7800	fmul qword ptr [0x789a18]
0060218d	d95df4	fstp dword ptr [ebp - 0xc]
00602190	85-0	test eax, eax
00602192	7e7a	jle 0x60220e
00602194	57	push edi
00602195	8b4508	mov eax, dword ptr [ebp + 8]
00602198	8d3c <mark>9</mark> 1	lea edi, dword ptr [eax + ebx*4]
0060219b	d907	fld dword ptr [edi]
0060219d	51	push ecx
0060219e	d87510	fdiv dword ptr [ebp + 0x10]
006021a1	d9 <mark>5d£8</mark>	fstp dword ptr [ebp - 8]
006021a4	d94518	fld dword ptr [ebp - 8]
006021a7	d91c24	fstp dword ptr [esp]
006021aa	eBb6feffff	call 0x602065
006021af	d95d0c	fstp dword ptr [ebp + 0xc]
006021b2	d9450c	fld dword ptr [ebp + 0xc]

- use disasm state as additional context
 - select different mixer weights, indirect probs, etc
- use disasm state to maintain different history buffers
 - one per state mostly, some for multiple states
 - represent "last opcode bytes", "last displacement bytes", etc
- double the number of context models in code section
 - \circ $\,$ second set only looks at history for current disasm state $\,$
 - 0 21 * 2 * 3 + 1 + 8 = 135 predictions!
 - \circ single wasted SIMD lane during mixing (136 / 8 = 17)
- best of both worlds:
 - \circ model can find correlations in in-place code
 - model can find correlations in "reordered" code (history buffers)

- possibility to mix/match histories in more arbitrary ways
 - \circ main motivation behind (so far failed) genetic algorithm idea
 - might still be interesting, needs more experimentation
- leads to larger model due to x86-specific stuff
- much can be folded into the main compressed data
 - as long as it's decompressed before code section
- additional model logic/history buffer code still somewhat big
- this is why a second compressor stage makes sense :)



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wrap-up

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- MAKE MORE INTROS!!!!!!!!!!

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- packers are fun!
- MAKE MORE INTROS!!!!!!!!!!
 - we have a synth that can help too... :)

Finished release [optimized] target(s) in 0.01s

Running `target\release\squishy.exe 'C:/Program Files (x86)/Steam/steamapps/common/DOOMEternal/DOOMEternalx64vk.exe' ./out.exe` squishy 0.1.0 | made with <3 by Jake "ferris" Taylor / logicoma 2016-2020

- big squish: 510674660 -> 16213766 (96.83%) in 888.75s (~561.13kb/s)

thread 'main' panicked at 'Compressed size too large; can't adjust image base to make room for compressed image.', src\main.rs:357:9 note: run with `RUST_BACKTRACE=1` environment variable to display a backtrace

error: process didn't exit successfully: `target\release\squishy.exe 'C:/Program Files (x86)/Steam/steamapps/common/DOOMEternal/DOOM Eternalx64vk.exe' ./out.exe` (exit code: 101)

many thank, wow jake "ferris" taylor / logicoma @ferristweetsnow yupferris at gee-mail

turn back, here be dragons

FORK ENDS HERE BRUH

- Asdfasdfasdfasdfasfasdf
- Asdf
- Asd
- Fa
- Sdf
- Asdf
- Asdf
- Asd
- Fasd
- f

- let's take an example, our string from earlier: AABC
- armed with a model: p = { A: 50%, B: 25%, C: 25% }
- create a new encoding by assigning new bit strings to the original symbols
- intuitively, make more common symbols use fewer bits than less common symbols

- let's take an example, our string from earlier: AABC
- armed with a model: p = { A: 50%, B: 25%, C: 25% }
- and an encoding: { A: 0, B: 10, C: 11 }
- encode our string with our encoding via per-symbol substitution

- A A B C
- 0 0 10 11 -> 001011
- the same string is 6 bits in our new representation!

- decode our string with our encoding via per-symbol substitution

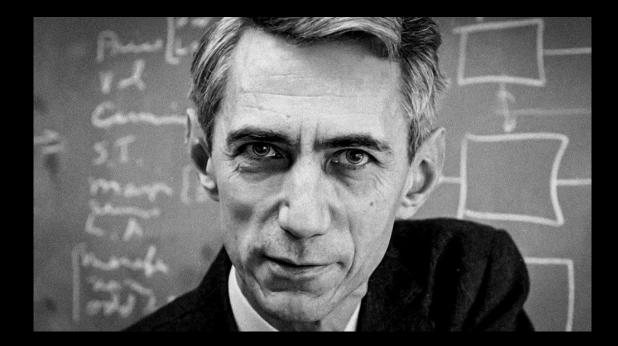
- 001011 -> 0 0 10 11
- A A B C
- it works!

- this was a contrived example, but you just learned a lot!
- our encoding was an example of a prefix code (just like huffman)
- it was also an optimal code for our model
 - this means that given the same model, we can not make a representation that would code the string using fewer bits than this!

• wait what how

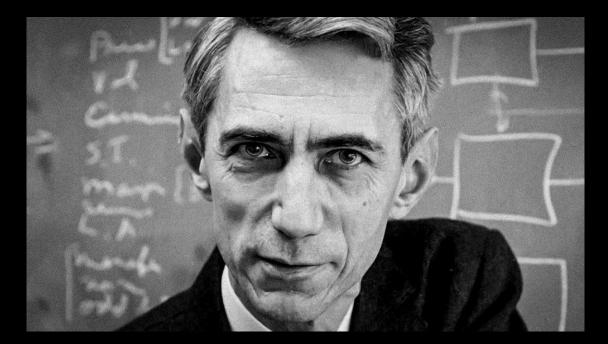


• meet claude shannon



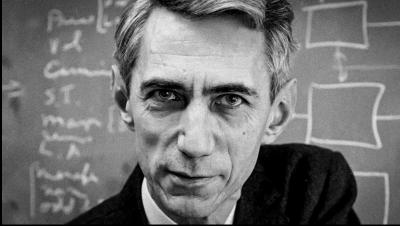
- meet claude shannon
- look at that hot piece of man(tropy)

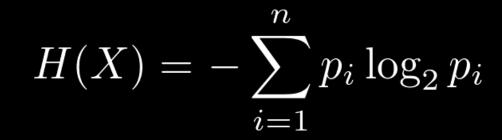
 ok this pic is haunting af but hear me out

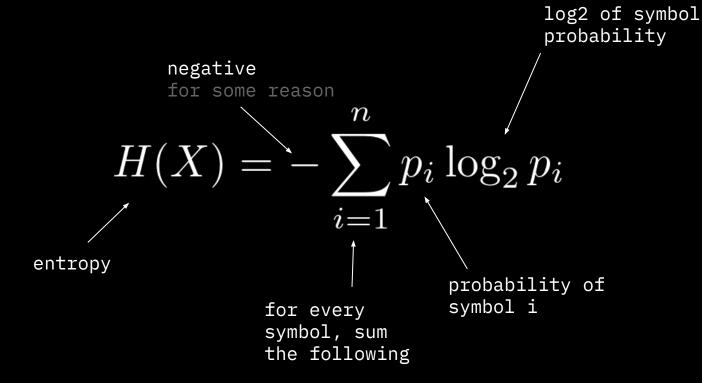


- claudy with a shans of meatballs over here did this really cool thing
 - he actually did a butt ton of awesome stuff!!!
- he came up with a way to quantify the average information content of a string
- he called it entropy
- and it goes a little somethin like this:

$$H(X) = -\sum_{i=1}^{n} p_i \log_2 p_i$$





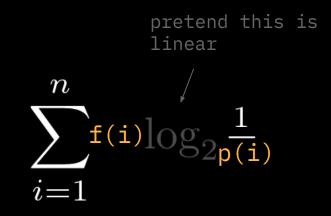


n $-\sum p_i \log_2 p_i$ $\overline{i=1}$

n $\sum p_i \log_2 \frac{1}{p_i}$ i=1

n $f(i)\log_{2p(i)}$ i=1

not exactly equivalent but it scales the same which is the important part so pls ignore



n

i=1

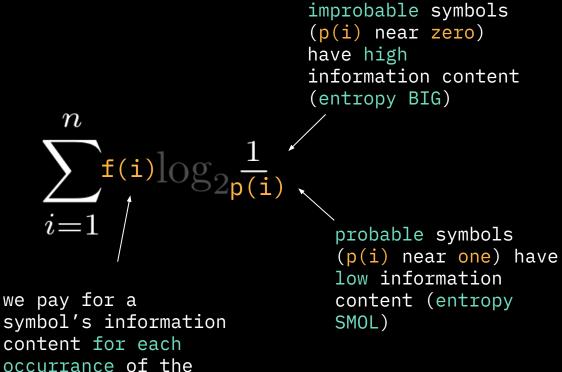
f(i)

p(i)

improbable symbols
(p(i) near zero)
have high
information content
(entropy BIG)

```
improbable symbols
                        (p(i) near zero)
                        have high
                        information content
                        (entropy BIG)
 n
    f(i)]0
                p(i)
i=1
                          probable symbols
                          (p(i) near one) have
                          low information
                          content (entropy
                          SMOL)
```

symbol!

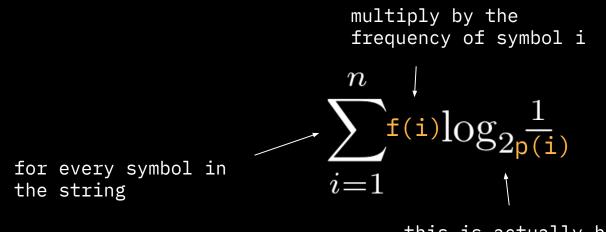


 $\sum_{i=1}^{n} f(i) \log_{2p(i)} \frac{1}{(i)}$

n $f(i)\log_{2p(i)}$

 $n \sum_{i=1}^{n} f(i) \log_{2p(i)} \frac{1}{i}$ this is actually how many bits symbol i should be coded with optimally!

multiply by the frequency of symbol i nf(i)]0g₂ (i) i=1this is actually how many bits symbol i should be coded with optimally!



this is actually how many bits symbol i should be coded with optimally!

n $\int f(i) \log_{2p} \frac{1}{(i)}$ i=1

this equation gives us the optimal number of bits we can use to code our string!!!

n $\int f(i) \log_{2p} \frac{1}{(i)}$ i=1



this equation gives us the optimal number of bits we can use to code our string!!!

n $-\sum p_i \log_2 p_i$ i=1

the real entropy equation is just a normalized version of that :)

• let's apply this to our example

optimal_bits(AABC) =
 optimal_bits(AA) +
 optimal_bits(B) +
 optimal_bits(C)

• let's apply this to our example

optimal_bits(AABC) =
 2 * log2(1 / 1/2) +
 1 * log2(1 / 1/4) +
 1 * log2(1 / 1/4)

• let's apply this to our example

• optimal_bits(AABC) =
 2 * log2(2) +
 1 * log2(4) +
 1 * log2(4)

• let's apply this to our example

- optimal_bits(AABC) =
 - 2 * 1 +
 - 1 * 2 +
 - 1 * 2

• let's apply this to our example

- optimal_bits(AABC) =
 - 2 +
 - 2 +
 - 2

• let's apply this to our example

• optimal_bits(AABC) = 6

• let's apply this to our example

• optimal_bits(AABC) = 6

• neato burrito