System Security

Introduction to Minix (Lecture 3)
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Outline of today

- Minix architecture
- Working with Minix
- Your first Minix Application
- Your first Minix Device Driver

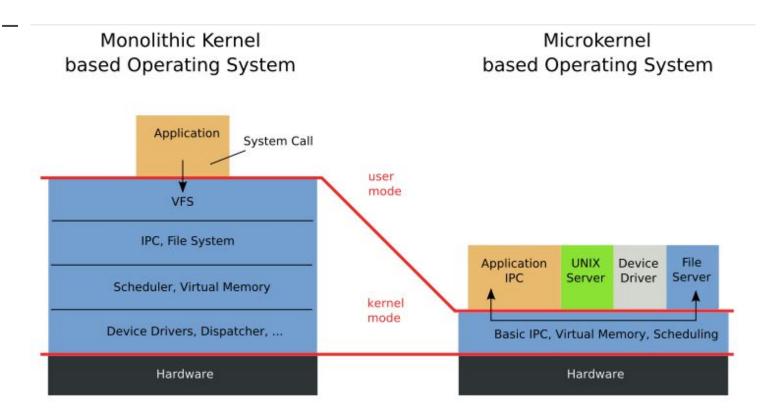
Minix Architexture

```
0 (NASA)/ 1 (MS)/ 15 (Avg)
Bugs per 1000 LOC
```

https://www.vidarholen.net/contents/wordcount/

About Minix

- Minix Mini Unix (Minix) basically, a UNIX compatible operating system.
- Minix is small in size, with microkernel-based design.
- Minix has been kept (relatively) small and simple.
- A POSIX-compliant, preemptive multitasking operating system
- Support for TCP/IP
- Source code: C language
- BSD licence
- Support for x86, ARMv7



User Processes

Device Drivers

Servers

Kernel

https://wiki.minix3.org/doku.php?id=developersguide:overviewofminixarchitecture

User Processes

Device Drivers

Servers

- central component of the OS
- manages the system's resources
 - o interface between hardware and software
- lowest-level of abstraction
- only component executed in kernel mode
 - Interrupts
 - Scheduling
 - Processes
 - Inter Process Communication (IPC)

User Processes

Device Drivers

Servers

- Offer services to the rest of the system
- Interact with drivers/servers/kernel
- Virtual Memory
 - Allocates physical memory
 - Changes page tables
- Virtual File System
 - Abstraction from of the actual file systems
 - Has special privileges to allows device drivers to access memory of other processes
- Actual File System
- Process Manager
 - Handles process informations
- Reincarnation Server
 - Start/ping/restart other servers
- Network stack

User Processes

Device Drivers

Servers

- handle a specific hardware component.
- each driver has access through the micro-kernel to its hardware resources (input/output ports, interrupts)
- mass storage (floppy disk, hard disk, CD/DVD/Blu-Ray/tape drive...)
- network card
- printer and serial port drivers.

User Processes

Device Drivers

Servers

- User-land processes are processes just like drivers and servers
- they do not have specific privileges
- to access facilities such as the file system or the network, they perform requests to servers
- application programming interface (API) is POSIX, a set of common system calls (not the Minix micro-kernel ones) and libraries functions all Unix-like operating systems share.
 - Mkdir, fork, execl, pipe, kill, networking, mmap

Reliability

https://wiki.minix3.org/doku.php?id=www:documentation:reliability

- 1. Reduce kernel size
- 2. Cage the bugs
- 3. Limit drivers' memory access
- 4. Survive bad pointers (reincarnation)
- 5. Tame infinite loops (reincarnation)
- 6. Limit damage from buffer overruns (partially true)
- 7. Restrict access to kernel functions
- 8. Restrict access to I/O ports
- 9. Restrict communication with OS components
- 10. Reincarnate dead or sick drivers
- 11. Integrate interrupts and messages

Message Passing

https://wiki.minix3.org/doku.php?id=developersguide:messagepassing

- fixed-length 64 bytes of data
 - The endpoint: a 4-byte identifier of who receives the message
 - Message type: a 4-byte message type identifier
 - Payload: 56 bytes of data
 - If more that 56 bytes, use memory grants
- each endpoint has a 1 message receive buffer inside the micro-kernel
- API
 - SEND: the sender is blocked until the message is delivered
 - RECEIVE: the process is blocked until a message is delivered to them
 - SENDREC: the sender is blocked until it receives a reply from the receiver.

Memory Grants

https://wiki.minix3.org/doku.php?id=developersguide:memorygrants

- allow processes to transfer large amounts of data
- processes grant access to regions of memory to another process
- attributes:
 - Access mode: read/write/read-write
 - Virtual memory region (byte granularity)
- cpf_* functions wrap system calls to manage grants
- types:
 - Direct
 - Indirect
 - Magic (only VFS)
- grant-id communicated via IPC
- grantee uses safecopies syscall specifying grant-id

Developing Minix

Developing for Minix

- Cross compiling:
 - building MINIX from another operating system
 - suggestion use Ubuntu 18.04
 - There is a VM, but is slow
- Native compiling:
 - building MINIX from within MINIX
 - High probability that you break everything
 - you change order of arguments of a system call, and you change accordingly the useland libs
 - You rebuild the system, the kernel and userland are recompiled and installed
 - You are still running the old kernel, but new processes (i.e. GCC) start to use the new userland

Crosscompiling

- install dependencies
 - a. apt-get install build-essential curl git zlibc zlib1g zlib1g-dev g++
- fetch sources
 - a. https://gits-15.sys.kth.se/robertog/system-sec-minix
- bash ./releasetools/x86_hdimage.sh
 - a. Build the toolchain (i.e. clang)
 - b. Build kernel, server, driver, libs, applications (i.e. vi), games
 - c. Package everything and create a disk image
- test with qemu (or your own computer if you are really brave)

Crosscompiling

- fetch sources
 - a. git clone git://git.minix3.org/minix
 - b. git clone https://gits-15.sys.kth.se/robertog/system-sec-minix.git
- use "master" (other branches are examples)

Crosscompiling

- bash ./releasetools/x86_hdimage.sh
 - a. Build the toolchain (i.e. clang)
 - First time takes several hours
 - There are problems to link LLVM, either fix the link script or compile adding -l termcap
 - It is recompiled if you update your Linux kernel
 - b. Build kernel, server, driver, libs, applications (i.e. vi), games
 - c. Package everything and create a disk image
- It is not needed to use release tool every time, but until you are confident it highly recommended
- To speed up set \${JOBS=3} in ./releasetools/image.defaults

https://wiki.minix3.org/doku.php?id=usersguide:runningonqemu

- ./releasetools/x86_hdimage.sh prints the qemu command to run the vm
- there is little to no support for external debuggers a. You can debug the microkernel, but difficult to debug processes
- Install KVM, it makes the emulation faster
- qemu-system-i386 --enable-kvm -m 256 -hda minix_x86.img

Demo Session

OEMU



--- Welcome to MINIX 3. This is the boot monitor. ---1emory: 639/129920 k

- 1. Start MINIX 3
- 2. Start latest MINIX 3
- 3. Start latest MINIX 3 in single user mode
- 4. Start MINIX 3 ALIX
- 5. Edit menu option
- 6. Drop to boot prompt
- 7. Start MINIX 3 (3.4.0)

Choose an option; RETURN for default; SPACE to stop countdown. Iption 2 will be chosen in <u>2</u> seconds.

You first Minix Program

"Hello world!"

Developing for Minix https://wiki.minix3.org/doku.php?id=developersguide:programmingminix

- Crosscompilation or In Minix compilation
- Good text about secure programming
 - https://dwheeler.com/secure-programs/Secure-Programs-HOWTO/index.html

The Minix Hello World

- games/demo01/Makefile
 - o BSD Makefile
 - https://wiki.netbsd.org/bsd make/
- games/Makefile
 - Add demo01 here
- distrib/sets/lists/games/mi and distrib/sets/lists/minix-games/mi
 - Add ./usr/games/demo01
- games/demo01/demo01.c
 - Your hello world app

The Minix Hello World

- Compile and deploy:
 - ./releasetools/x86_hdimage.sh
- Only compile:
 - o PATH=\$PATH:<???>/obj.i386/tooldir.Linux-<???>/bin/
 - o nbmake-i386 clean
 - o nbmake-i386

Your first Device Driver

User Processes

Device Drivers

Servers

- handle a specific hardware component.
- each driver has access through the micro-kernel to its hardware resources (input/output ports, interrupts)
- mass storage (floppy disk, hard disk, CD/DVD/Blu-Ray/tape drive...)
- network card
- printer and serial port drivers.

Driver Architecture

- The drivers source code are located in minix/drivers
- The binary executables are located under the /service/ directory

Character Based

- data as a stream of bytes
- do not support the seek
- Printer / Serial Port / TTY / Random/

Block Based

- data in chunks called blocks of a constant number of bytes
- Support seek
 - access to any block at any moment
- hard disk and CD drives

Driver Architecture

- Interact with Kernel/VFS/Reincarnation server using specialized Low level protocol
 - Using message passing, memory grants and system calls
- High-level abstraction libraries
 - o provide a higher-level abstraction interface
 - o shield the drivers from low-level protocols changes

- blockdriver: for block-oriented drivers
- chardriver: for generic character-oriented drivers
- inputdriver: for drivers communicating with the input server
- netdriver: for network device drivers

A simple character based device

- Driver source code
 - minix/drivers/mydriver/mydriver.c
 - o minix/drivers/mydriver/mydriver.h
- Configuration:
 - o minix/drivers/mydriver/mydriver.conf
 - o Allows driver to communicate to other services and invoke syscalls
- Makefile
 - minix/drivers/mydriver/Makefile
 - o It copies the driver binary in /service and conf in /etc/system.conf.d
 - minix/drivers/Makefile
 - Add mydriver here
- Distribution List
 - distrib/sets/lists/minix-base/mi
 - add /etc/system.conf.d/mydriver and /service/mydriver

The System Event Framework (SEF)

- System library to deal with system events
- sef_startup()
 - RS will send the driver an initialization message with information on how to initialize properly
 - The driver is expected to reply back with OK or error
- 1. Driver registers callback
- Driver calls sef_startup()
- 3. SEF do the rest and invokes callbacks

A simple character based device

- To start the driver
 - o minix-service up /service/mydriver
- To stop the driver
 - minix-service down mydriver
- Reincarnation Server (RS)
 - Start and stop services
 - minix-service sends a message to RS
 - RS pings the driver
 - If the driver aborted or exited
 - Restart the driver
 - If the driver never replies since is is stuck in a loop
 - Stop and restart
- Driver should reply to keep-a-live messages
 - o library libchardriver takes care of various common tasks

The System Event Framework (SEF)

- sef_setcb_init_fresh
 - fresh start initialization
- sef_setcb_init_lu
 - o init after live update
 - there are specific call backs to inform RS that we can be updated and prepare for update
- sef_setcb_init_restart
 - o init after restart

The Chardriver

- struct chardriver
 - Call backs for char based device
 - chardriver library has a main-loop to fetch messages, invoke callbacks, send reply messages
- open/close/read/write
- sys_safecopyto(endpt, grant, 0, (vir_bytes) ptr, size)
 - endpt: id of process requesting read (actually is the process that granted the memory access)
 - o grant: grant to a region of memory of another process
 - 0: offset in the receiver buffer
 - ptr: address in driver's virtual memory
 - size: amount of bytes
 - o actual copy done by the kernel

The Chardriver final notes

- minix/include/minix/dmap.h
 - o Allocate major 18 to mydriver
- minix/commands/MAKEDEV/MAKEDEV.sh
 - Make device file /dev/mydriver at boot
- etc/usr/rc
 - Start driver at boot

The Chardriver

- > mknod /dev/mydriver c 18 0
 - major ID 18 (look at minix/dmap.h). Identifies device driver
 - character device with major number 18 and minor number 0
- > minix-service up /service/mydriver -dev /dev/mydriver
- > cat /dev/mydriver
- > minix-service refresh mydriver
- > minix-service down mydriver

Example projects

- Anti-code injection (2018)
 - Monitor exec to only execute signed code / prevent bad programs
 - Inspect process text area to detect alterations / check signatures
- Control flow integrity (2018)
 - Monitor syscalls and enforce caller PC integrity
- Firewall (2018)
- Run-time patching (2018)
- Malware monitor and analysis
- Anti DOS
 - Monitor memory allocation / fork / fopen to limit used resources
- Anti buffer overflows
 - Inspect stack to identify stack corruptions
- Extend GDB support
- Re-randomization

Questions?