Mempool as a Battleground:

RBF Pinning, package relay, v3, ephemeral anchors

glozow

Today

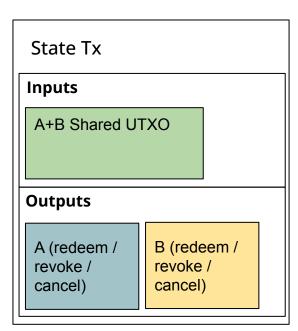
- The Problem
- Current Options
- Solution Part 1: Package Relay
- Solution Part 2: v3
- Solution Part 3: Ephemeral Anchors

The Problem

L2 == awesome

Sign now, broadcast later.

- do more stuff, put less on-chain
- privacyTM, scalabilityTM



L2 == awesome

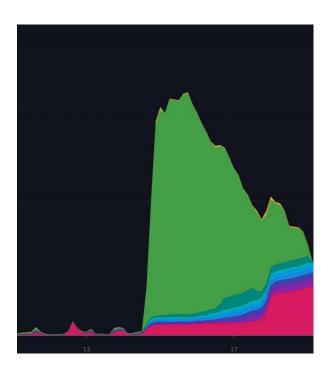
Sign now, broadcast later.

- do more stuff, put less on-chain
- privacyTM, scalabilityTM

The problem

Typically, when you sign, you decide fees.

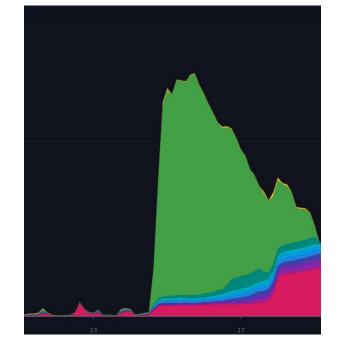
- A lot can change between sign and broadcast.
- This tx is shared with someone untrusted.



Current Options

Predict the Fee Using Your Magic Crystal Ball

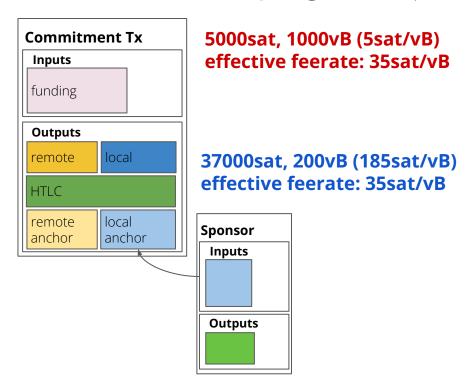




Related methods:

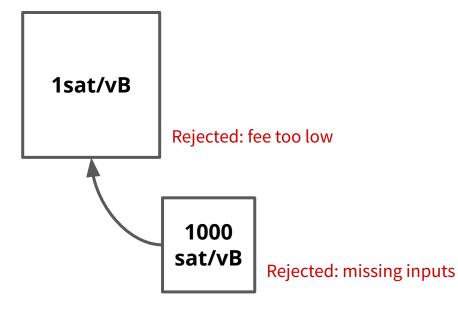
- Overshoot the feerate
- Sign multiple transactions at different feerates (Bastien Teinturier: https://lists.linuxfoundation.org/pipermail/lightning-dev/2022-October/003729.html)

Attach a fee-bumping child (CPFP)



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Caveat: parent must meet mempool min feerate



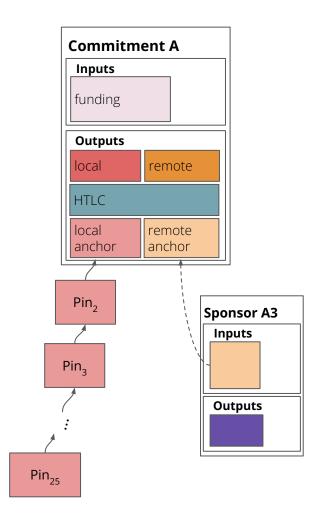


mempool.space/ 2023-03-08

Attach a fee-bumping child (CPFP)

Ugliness of anchor outputs:

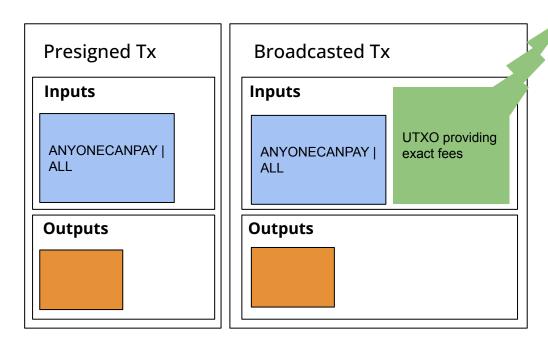
- Needs one for each participant
- Various hacks needed to avoid pinning
 - Other outputs can't be spent (CSV 1)
 - Needs CPFP carve out (2-party only)
- Shaved off from channel balance (cannot be dust)
 - blocker for eltoo
 - Creates low-value UTXOs (unless cleaned up)



ANYONECANPAY to increase inputs

Sign transaction with SIGHASH_ANYONECANPAY

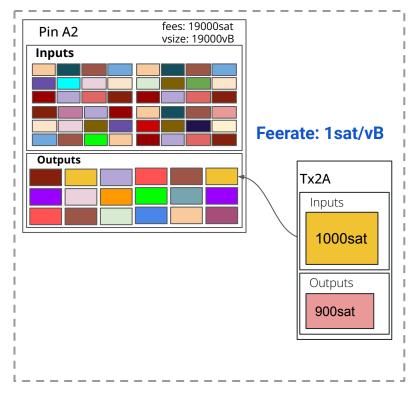
Adjust fees by adding inputs

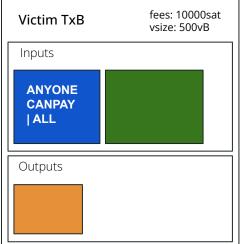


(Revault: https://github.com/revault/practical-revault/blob/master/transactions.md#cancel_tx)

ANYONECANPAY -> anyone can RBF

Attacker's transactions

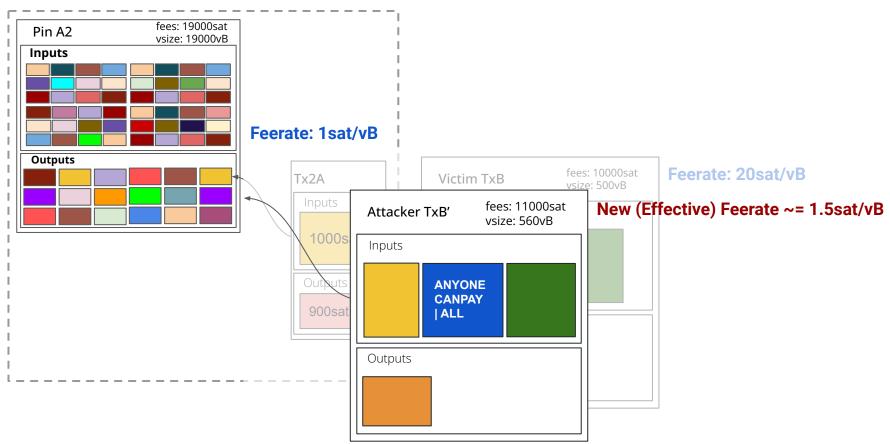




Feerate: 20sat/vB

ANYONECANPAY -> anyone can RBF

Attacker's transactions





mined or entering a mempool

advantage of mempool policy limitations to prevent a tx from getting

Pinning Attack: a type of censorship in which attacker takes

Pinning Attack: a type of censorship in which attacker takes advantage of mempool policy limitations to prevent a tx from getting mined or entering a mempool

attacker isn't paying fair price

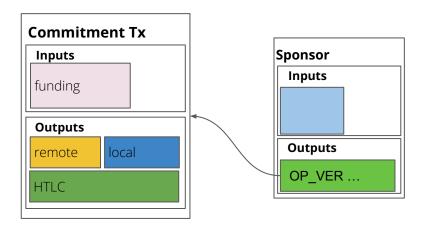
mempool should provide a fair (fee-based) market for block space

"Can we Soft Fork it out?"



Similar to CPFP: Transaction Sponsors soft fork

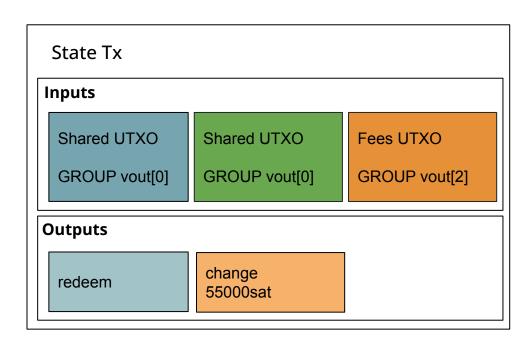
- no anchor outputs
- "anyone can bump"
- many similar limitations to CPFP
 - everything still needs to be CSV 1
 - package limit pinning
 (sponsor-sponsee ~= parent-child)
- needs soft fork





SIGHASH_GROUP / Signature bundles soft fork

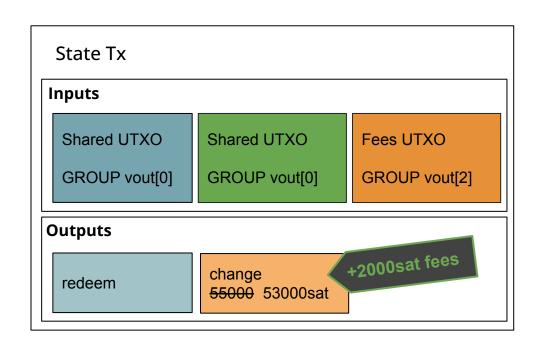
Instead of signing all/single/none of the outputs, specify a range



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Instead of signing all/single/none of the outputs, specify a range

Adjust fees simply by modifying change output amount

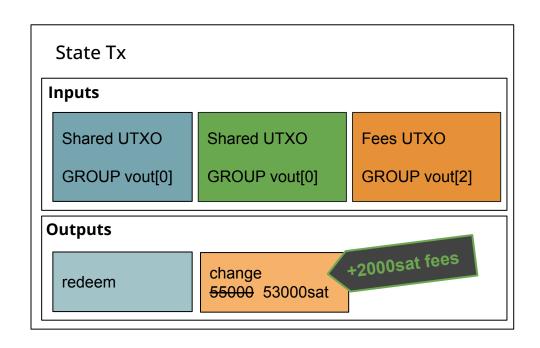


Anthony Towns: https://lists.linuxfoundation.org/pipermail/bitcoin-dev/2021-July/019243.html Rusty Russell: https://lists.linuxfoundation.org/pipermail/bitcoin-dev/2021-July/019243.html

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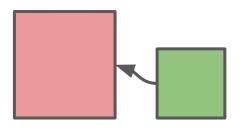


Solutions Categorized



Broadcast As Is

- Magical fee prediction
- Sign multiple txns



Add Sponsor Tx

- CPFP
- Transaction
 Sponsors (soft fork)

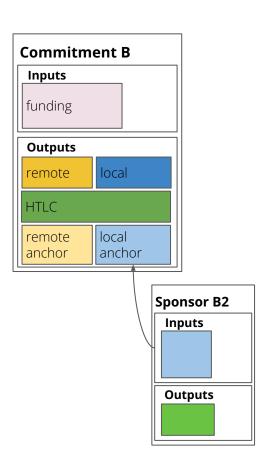


Modify the Tx Itself

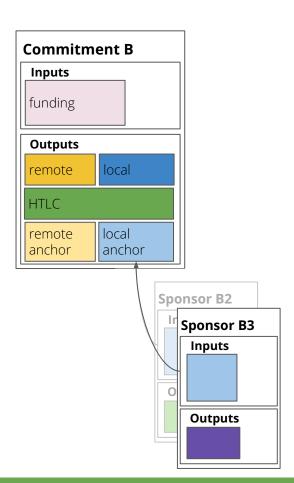
- ANYONECANPAY
- SIGHASH_GROUP (soft fork)

Solution Part 1: Package {CPFP, RBF, Relay}

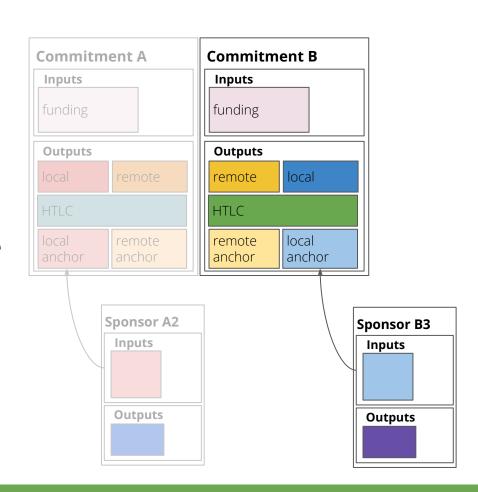
- 0 fees or 1sat/vB on shared tx (commitment)
- 🔽 add fees at broadcast time



- 0 fee or 1sat/vB on shared tx (commitment)
- add fees at broadcast time
- ✓ bump feerate by RBFing the child
- package relay protocol changes make propagation more reliable

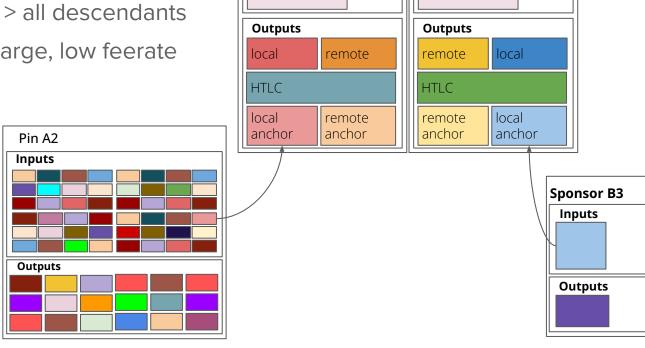


- 0 fees on shared tx (commitment)
- ✓ add fees at broadcast time
- ✓ bump feerate by RBFing the child
- package relay protocol changes make propagation more reliable
- if conflicting tx exists, child fees count for RBF fee-related rules



X Caveat: "Rule 3" Pinning

replacement fees must > all descendants descendant(s) may be large, low feerate



Commitment A

Inputs

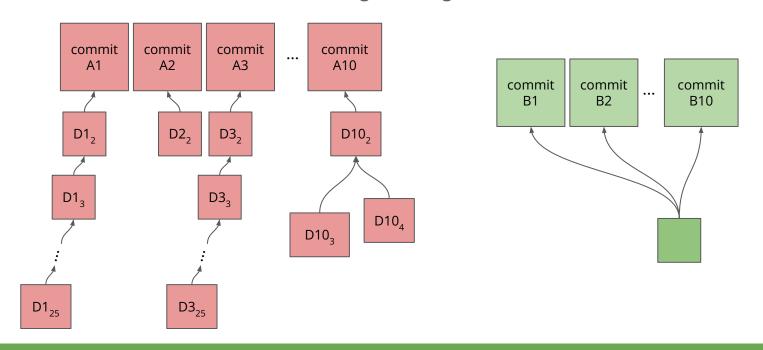
funding

Commitment B

Inputs

funding

X Caveat: "Rule 5" Pinning can't replace more than 100 at a time. batching is dangerous



"Ah ok, we just need fix RBF"

some idiot, January 2022



Solution for this: Let's add an incentive compatibility rule to RBF!

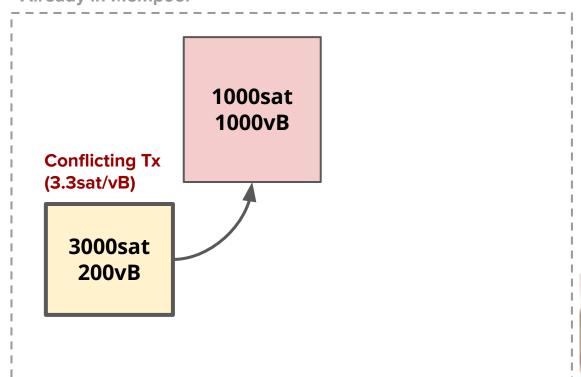
Already in Mempool

Conflicting Tx (15sat/vB)

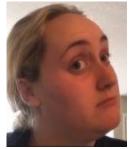
3000sat 200vB Replacement Tx (10sat/vB)



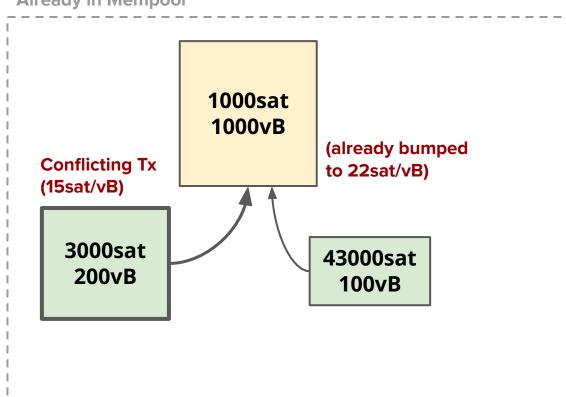
Already in Mempool



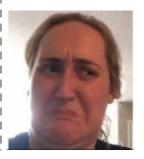
Replacement Tx (10sat/vB)



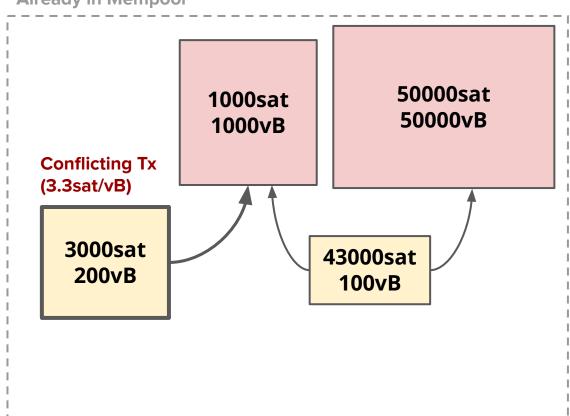
Already in Mempool



Replacement Tx (10sat/vB)



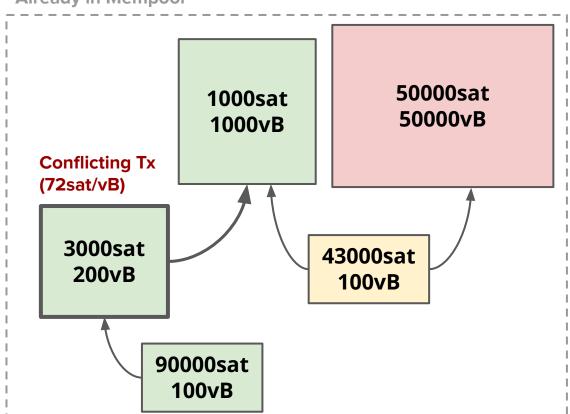
Already in Mempool



Replacement Tx (10sat/vB)



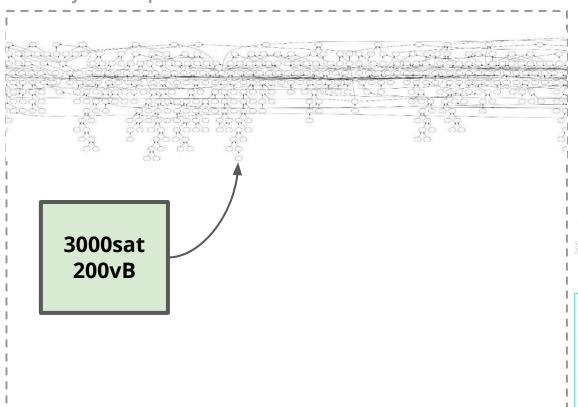
Already in Mempool



Replacement Tx (10sat/vB)



Already in Mempool



Replacement Tx (10sat/vB)

1000sat 100vB



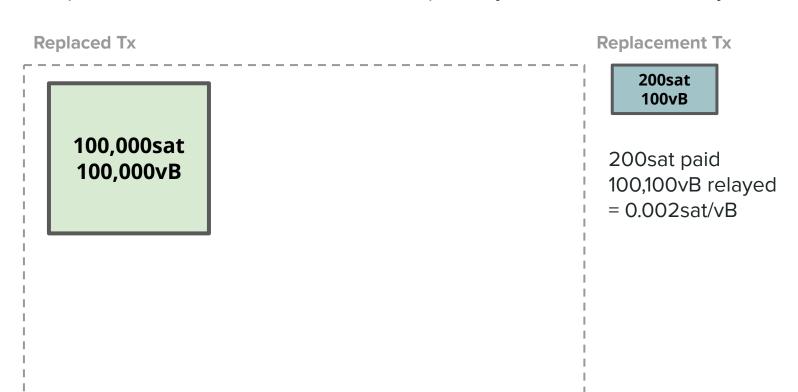


The above cluster was composed of 219 unconfirmed txs. I just found another cluster of 881 linked unconfirmed txs. 😝

If you work on coin selection, please take an input's full ancestry into account when evaluating the viability of unconfirmed inputs during transaction building.

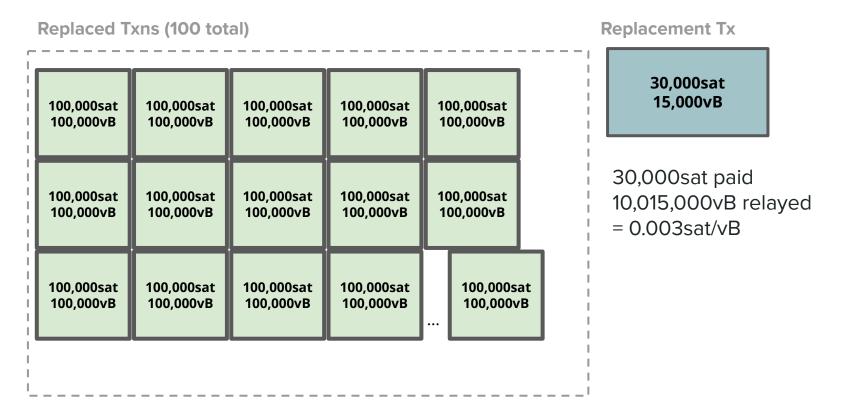
Before you say "can we get rid of Rule 3 entirely?"

"Replacement's feerate and incentive compatibility score must increase by 2x"



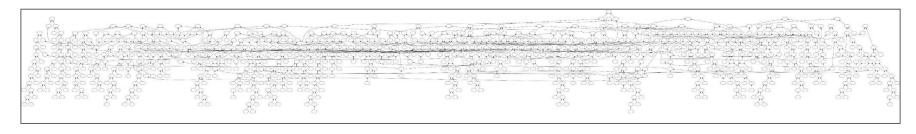
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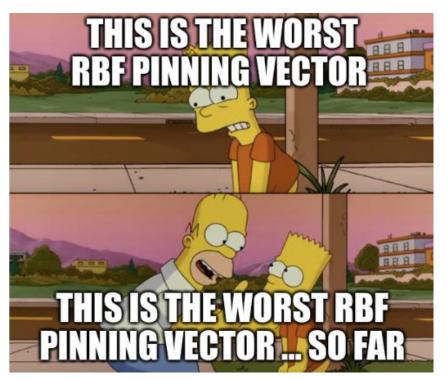
"Replacement's feerate and incentive compatibility score must increase by 2x"



Solution Part 2: v3 to fix pinning

Takeaway: we allow these, even though we can't properly handle them.

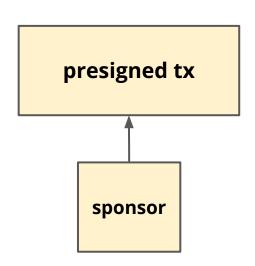




@instagibbs

V3 Rules

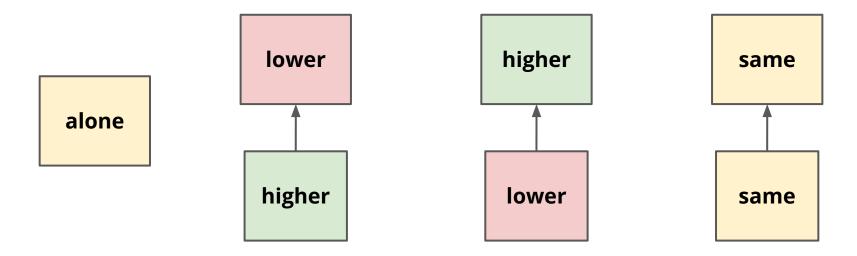
- 1 parent 1 child only
- child can't be more than 1000vB
- (unconfirmed) v3 must spend v3
- (unconfirmed) non-v3 must spend non-v3
- v3 signals replaceability



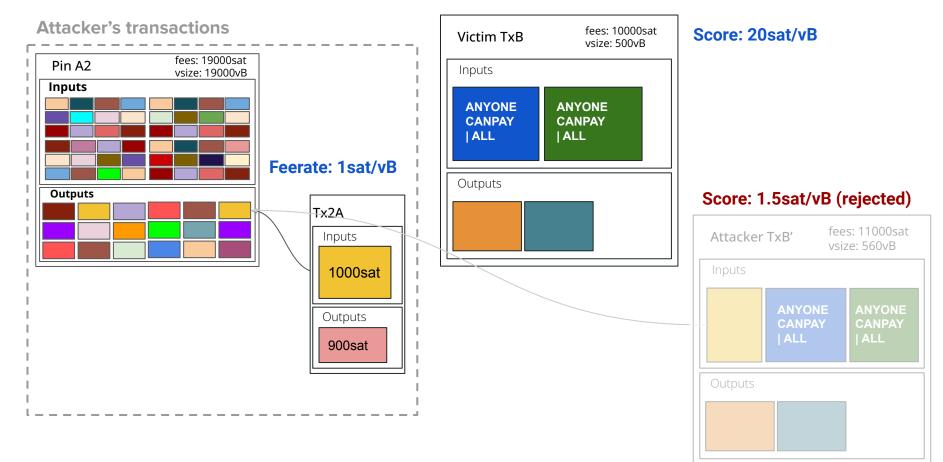
Incentive compatibility score

Cluster can't be larger than 2, so it's just = min(self feerate, ancestor feerate)

Pretty easy to show this is always correct:



ANYONECANPAY Replacement Pinning

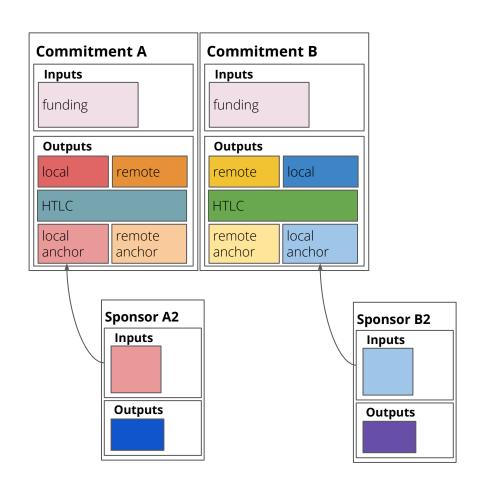


Rule 3 Pinning

Limiting the size of an attached tx == limiting the economic damage your counterparty can do to you

Need A or B to confirm?

- 1. Decide confirmation target, feerate is f
- 2. Assuming both txns have size \mathbf{s} vB, add fees of \mathbf{f}^* (\mathbf{s} + 1000) to fee-bumping child
- 3. Broadcast commitment tx + child
- 4. If no confirmation, must be because feerate too low. RBF the child



"Cute, but is this incentive compatible?"

Miner Benefits

(if users use it)

- DoS-resistant, generally computationally cheap to handle
- Can assess incentive compatibility quite easily

User Benefits

(if network adopts it)

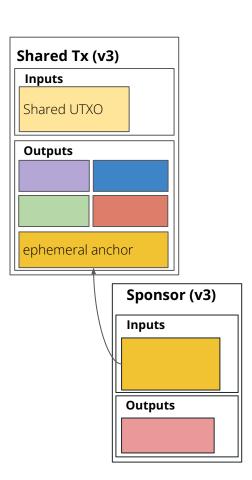
- No difference between tx signaling and its ancestor signaling
- Any RBF requires an incentive compatibility score increase
- ✓ Just broadcast, no need to monitor mempools to see if you need to pay extra to RBF
 - Rule 5 pinning severity reduced by 24x
 - Rule 3 pinning severity reduced by 100x

Solution Part 3: Ephemeral Anchors

(Greg Sanders)

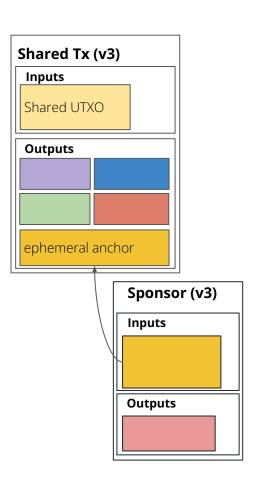
Ephemeral Anchor Rules

- parent:
 - 0 fee, so it must be bumped
 - 1 OP_TRUE output to attach fee-bumping child ("anchor")
 - anchor output **can be any value** (including 0)
 - v3: only 1 child allowed
- child:
 - spends the anchor ("ephemeral")
 - v3: only 1 parent allowed



Ephemeral, 0-value, Anchor Outputs

- Anyone can bump the tx
 - Watchtowers don't need keys
 - **✓** Works for transactions shared between N>2 parties
- Exactly 1 anchor output. That output *must* be spent.
 - ✓ Smaller tx size
 - CPFP carveout can be phased out
 - ✓ Don't need 1 CSV for the other outputs
- ✓ No need to shave value off channel balance (wen eltoo?)



Thanks!