

Part of the Carbon Language, under the [Apache License v2.0 with LLVM Exceptions](#).
SPDX-License-Identifier: [Apache-2.0](#) WITH [LLVM-exception](#)

Carbon Language - <http://github.com/carbon-language>

Multi-Thread Examples

Authors: danakj, ...

Status: Draft ▾

Created: 2025-10-24

Docs stored in [Carbon's Shared Drive](#)
[Access instructions](#)

Problem

Examples that demonstrate parallelism and concurrency needs for Carbon. Either code that should compile, or problematic code that should not.

Examples

Thread spawn

- No join, so no way to know when it's done in the caller.
- Passing in a pointer that is invalidated

None

```
fn F() {  
  var i: i32 = 0;  
  var p: const i32* = &i;  
  spawn(fn [p] => G(p));  
  // Bad: i is destroyed, p is invalidated but in use on another thread.  
}
```

- Holding a pointer that is invalidated, moving its owner

None

```
fn F() {  
  var i: Box(i32) = .Make(0);  
  var p: const i32* = i.Get();  
  spawn(fn [~i] => ~i);  
}
```

```

    // Bad: i may be destroyed, then p is invalidated.
    Core.Print(*p);
}

```

- Holding a pointer that is invalidated, not moving its owner

```

None
var i: Cell(Optional(i32)): .Make(.Some(0));

fn F() {
    var p: const i32* = i.Get();
    spawn(fn => i.Set(.None));
    // Bad: The i32 in i may be destroyed/inactive, then p is invalidated.
    Core.Print(*p);
}

```

- Long lived vs short lived pointers

```

None

fn F(pointer_to_string_literal: const char*, pointer_to_string: const char*) {
    spawn(fn [let] => H(pointer_to_string_literal, pointer_to_string));
}

fn G1() {
    let world = "world" as strbuf;
    F("hello", world.Get());
    // Bad: `world` is freed, but the pointer is in use on another thread.
}

fn G2() {
    F("hello", "world");
    // OK: neither pointer is freed while in use on another thread.
}

```

- Long lived with mutation

None

```
fn F(pointer_to_global: const i32*) {
    spawn(fn [let] => H(pointer_to_global));
}

var g: i32 = 2;

fn G() {
    let world = "world" as strbuf;
    F(&g);
    // Bad: Writing to `g: i32` on this thread, but maybe reading it on another
    thread.
    g = 3;
}
```

- Passing unowned along with ownership

None

```
fn F() {
    var b: Box(i32) = .Make(2);
    var p: i32* = b.Get();
    spawn(fn [~b, p] => H(~b, p));
    // Bad: `p` needs to be invalidated for soundness. The box may be destroyed
    // at any time.
    Core.Print(*p);
}
```

- Passing combined owned, unowned, and global/static, pointer to locally-owned, and Rc.

None

```
class C(^B, ^G, ^L) {
    // Can move across threads.
    var b: Box(i32);
    // Moves with the owner, is valid. Must invalidate any copies on the original
    thread.
    var p: i32 ^B *;
    // Global so can move across threads. Must not invalidate any copies on the
    original
    // thread.
```

```

var g: i32 ^G *;
// Moves across threads without the other, not valid. Must be an error.
var l: i32 ^L *;
// Owns its data, but can not move across threads. Must be an error.
var r: Rc(i32);
}

var globalint: i32 = 3;

fn F() {
  var b: Box(i32) = .Make(2);
  var p: i32* = &c.b;
  var g: i32* = &globalint;
  var l: i32 = 1;
  var r: Rc(i32) = .Make(4);
  var c: C = { .b = ~b, .p = p, .g = g, .l = l, .r = r };
  // Two errors must occur:
  // - Sending `.l` across threads without its owner. The owner will invalidate
  //   it in the future.
  // - Sending `.r` across threads, as its _type_ is not valid to move across
  //   threads even though it owns its data.
  spawn(fn [~c] => H(~c));
  // Error: `p` needs to be invalidated for soundness. The box may be destroyed
  // at any time.
  Core.Print(*p);
  // Accepted: `g` must *not* be invalidated, as it points to a global not an
  // owned thing.
  Core.Print(*g);
}

```

Send effect

Maybe we need a **send** effect/function capability. Like we can say **invalidate(x)** we could say **send(x)** to say we move it (or a copy maybe?) to another thread. Then we can say no based on conditions like:

- **Rc** is **not** allowed (how do we mark it as such? Type capabilities?)
- **.l** is allowed by type but then we can look to see there's **not** an owner going with? And see there is not.
- **.b** is allowed by type and we can see that there **is** an owner going with.
- **send(x)** would invalidate any alias sets owned by something in **x**. So **p** gets invalidated, but **g** doesn't.

None

```
fn spawn[Func: Call](func: Func) [[send(^func.any)]];
```

- Problem is that anything that gets `send` applied to it needs to pass that up the call stack to anything that passed it along. Rust only requires putting `Send` on generics that pass things through a call stack, since types get `Send` as an auto trait.

None

```
class C;
```

```
fn GenericSpawns[T: ! type](c: T) [[send(c)]];
fn SpawnsGenerically(c: C) [[send(c)]];
fn CallsSpawnsGenerically(c: C) [[send(c)]];
```

Vs Rust

None

```
struct C;
```

```
fn GenericSpawns<T: Send>(c: T);
fn SpawnsGenerically(c: C);
fn CallsSpawnsGenerically(c: C);
```

Scoped thread spawn

- Mostly for the ability to pass in pointers to data in the caller, knowing when the thread is done with them.
 - To allow **immutable** sharing
 - Not to allow destruction from the thread
- A fixed point in the caller where the task is done running.
- Either scoped as a lambda: gives a nested scope for the point where the thread ends
- Or as a return guard that *must be destroyed*: allows returning the guard, giving a multi-function/non-local scope where the thread may be running
- `spawn_scoped` allows pointers to things before the scope

None

```
fn F() {
    var i: i32 = 0;
```

```

var p: const i32* = &i;
var scope: auto = spawn_scoped(fn [p] => G(p));
~scope;
// Scope does not finish destruction until the scoped_spawn is joined:
// i was not destroyed and p was not invalidated.
}

```

- But not if they are mutated in the scope.

```

None
fn F() {
  var i: Optional(i32) = 0;
  var p: const i32* = &i.Get();
  var scope: auto = spawn_scoped(fn [p] => G(p));
  // Bad, the i32 in i is gone, so p is invalidated but it's in use on another
  thread.
  i = .None;
  ~scope;
}

```

- And doesn't allow holding pointers to things given to the thread

```

None
fn F() {
  var i: Box(i32) = .Make(0);
  var p: const i32* = i.Get();
  var scope: auto = spawn_scoped(fn [~i] => ~i);
  // Still bad: i may be destroyed, so p may be invalidated.
  Core.Print(*p);
  ~scope;
}

```

Coroutines

- Web server calling backends

```

None
class Request {
  var data: buf(u8),
}
class Headers {
  var mode: Optional(i32);
}

```

```

    var pos: slice(u8); // In Request.data.
}

/*async*/ fn ParseHeaders(let r: Request) -> Future(Headers);
/*async*/ fn QueryDatabase(let h: Headers) -> Future(strbuf);
/*async*/ fn SendReply(reply: strbuf) -> Future(());

/*async*/ fn HandleRequest(var r: Request) -> Future<()> {
    let headers: Headers = ParseHeaders(r).await;
    var answer: strbuf = QueryDatabase(headers).await;
    return SendReply(answer); // returns the Promise/Future instead of .await.
}

/*async*/ fn ClearsDataDuringHandleRequest(var r: Request) -> Future<()> {
    let headers: Headers = ParseHeaders(r).await;
    // Bad: Invalidates headers.pos, but headers is used below.
    r.data.clear();
    var answer: strbuf = QueryDatabase(headers).await;
    return SendReply(answer); // returns the Promise/Future instead of .await.
}

/*async*/ fn RequestDoesntLiveLongEnough(let r: Request*) -> Future<()> {
    let headers: Headers = ParseHeaders(*r).await;
    // Bad: r could be invalidated as soon as we call .await, while ParseHeaders()
    // is working if it has any .await inside.
    return SendReply("done"); // returns the Promise/Future instead of .await.
}

```

Fiber race

(adapted from a Rust example)

```

None
// Run each of the supplied children concurrently on its own
fiber
// that is a child of the parent making this call. Return the
result
// of first child to finish, cancelling and joining the

```

```

// others.
//
// In Rust:
// pub fn race<'a, T>(children: Vec<Box<dyn 'a + FnOnce() -> T +
// Send>>) -> T

fn race[T: ! Type](children: Vec<Box<call!() -> T>> -> T {
    // There is no correct answer if the set of children is
    empty.
    assert!(!children.is_empty());

    // Spawn child fibers that race to store their result into
    // a OnceLock. After one of them wins, cancel and join them
    all.
    // OnceLock is mutable, but shared in Rust
    // (using interior mutability)
    let once = OnceLock::Make();
    nursery::with_nursery([ref once](n: Nursery) {
        // Start children competing to be the first.
        let fibers: Vec<Fiber> = children
            .map(|child| {
                // Run child in a a fiber
                n.spawn(|| {
                    let _ = once.set(child());
                })
            })
            .collect();

        // Wait for one to win.
        once.wait();

        // We've got what we need. Clean up.
        for (f in fibers) {
            f.cancel();
            f.join();
        }
    })
}

```

```
});  
  
// By now our result is in the OnceLock.  
Box::Take(once.consume())  
}
```