

(Note: this is written for norm floating point numbers. For denorm, the exponent will always be $\text{exp} - \text{bias} + 1$, and we do NOT add $2^{\text{exp} - \text{bias} + 1}$ in the mantissa.)

Floating point to decimal

Assuming 1 sign bit, 8 exponent bits, 23 mantissa bits, and a bias of -127.

Example: 0x70DA0000

1. Convert to binary

0x70DA0000 \rightarrow **0 11100001 1011010000...**

2. Find the exponent

11100001 \rightarrow 225, so $225 - 127$ (bias) = **98**

3. Convert mantissa

Write out the mantissa bits with numbers underneath them. These numbers will start from the exponent we just calculated minus one, so $98 - 1 = 97$. Like this:

1	0	1	1	0	1	0	0	...	0
97	96	95	94	93	92	91	90	...	75

Now, every bottom number that has a 1 on top will be included in our final answer:

$$2^{98} + 2^{97} + 2^{95} + 2^{94} + 2^{92}$$

Our original exponent of 2^{98} will always be included in this summation as well (but NOT for denorm).

4. Don't forget the sign bit

If the sign bit is 1, don't forget to multiply by -1. In this case our sign bit is 0, so our number is positive, so we don't have to do anything here.

Final answer: $2^{98} + 2^{97} + 2^{95} + 2^{94} + 2^{92}$

Decimal to floating point

Assuming 1 sign bit, 8 exponent bits, 23 mantissa bits, and a bias of -127.

Example: 9.75

1. Find sign bit

In this case, **0**, since we have a positive number.

2. Write out number in exponents of 2

$$9.75 = 2^3 + 2^0 + 2^{-1} + 2^{-2}$$

3. Find exponent bits

Our leading exponent will be the exponent we want to represent here. In this case, our leading exponent (the biggest one) is 3. So we do, $3 + 127$ (bias) = 130 \rightarrow **10000010**

4. Find mantissa bits

This process is similar to what we did in step 3 from converting floating point to decimal. Starting from the leading exponent minus one, so $3 - 1 = 2$, if the exponent is present in our summation, we write a 1, otherwise we write a 0. Again, this is better seen in an example, so here it is. The bottom numbers are the exponents, the top are the mantissa bits.

Our original summation: $2^3 + 2^0 + 2^{-1} + 2^{-2}$

0	0	1	1	1	0	0	...	0
2	1	0	-1	-2	-3	-4	...	-20

5. Convert to hex

0 10000010 001110000... \rightarrow **0x411C0000**

Final answer: 0x411C0000

Intuition

So how does this all work? Well let's first forget about floating point for a sec.

If I asked you in a math class to do $100 * 1.305$, you would probably just multiply them. But, if you wanted to, you could do something like

$$\begin{aligned}10^2 * (1 * 10^0 + 3 * 10^{-1} + 5 * 10^{-3}) &= \mathbf{10^2} + \mathbf{3 * 10^1} + \mathbf{5 * 10^{-1}} \\ &= 130.5\end{aligned}$$

It's the same idea with floating point, but in this case thinking about it using the second method is easier (for me at least). Since we're in base 2 now, we replace the 10's with 2's.

So $1000 * 1.1001100\dots_2$ in decimal would be

$$\begin{aligned}2^3 * (1 * 2^0 + 1 * 2^{-1} + 1 * 2^{-4} + 1 * 2^{-5}) &= 2^3 + 1 * 2^2 + 1 * 2^{-1} + 1 * 2^{-2} \\ &= \mathbf{2^3} + \mathbf{2^2} + \mathbf{2^{-1}} + \mathbf{2^{-2}} \\ &= 12.75\end{aligned}$$

Helpful Tips

So besides conversion, how is this useful? Well, for example, what if I asked you for the smallest integer unrepresentable by the standard 1 sign, 8 exponent, 23 mantissa bit floating point representation? You've probably already seen this problem, but this method gives an intuitive way to think about it.

Consider the example from the first page with $2^{98} + 2^{97} + 2^{95} + 2^{94} + 2^{92}$. Could I represent $2^{98} + 1$? Nope, because $2^{98} + 2^0$ would be impossible to encode. (We know that to increase a floating point number by the smallest amount possible, we add 1 to the last bit, so...) Since there are only 23 mantissa bits, the lowest possible increase (by adding 1 to the last bit) would be 2^{75} . (Because in this case, a '1' in the last bit corresponds to adding 2^{75} to our answer!)

So to find the smallest integer unrepresentable, we just need to find when the last bit corresponds to adding 2^1 rather than 2^0 . (Adding 2^0 means adding 1, which hits every integer, but adding 2^1 means adding 2, which skips every other integer.) When the last bit corresponds to 2^1 , our exponent is 2^{24} . Therefore, the smallest integer unrepresentable is $2^{24} + 1$. (Try

representing $2^{24} + 1$ using this method, and you'll see why it's impossible. If my explanations make no sense, I'm sorry, but pls try doing this out, I think it'll help a lot.)

All in all, hopefully the logic makes sense, and you can try applying it to other floating point questions. This is how I do every floating point question :)

good luck!