## CSE 414 Section 7

## **Locking Worksheet**

- (a) Consider a concurrency control manager that uses strict two phase locking that schedules three transactions:
  - T1: R1(A), R1(B), W1(A), W1(B), Co1
  - T2: R2(B), W2(B), R2(C), W2(C), Co2
  - T3: R3(C), W3(C), R3(A), W3(A), Co3

Each transaction begins with its first read operation, and commits with the *Co* statement. Answer the following questions for each of the schedules below:

- Is this schedule possible under a strict 2PL protocol?
- If strict 2PL does not allow this schedule because it denies a read or write request, is the system in a deadlock at the time when the request is denied?

## i. Schedule 1:

```
R2(B), W2(B), R3(C), W3(C), R3(A), W3(A), Co3, R2(C), W2(C), Co2, R1(A), R1(B), W1(A), W1(B), Co1
```

- a) Is it possible under strict 2PL?
- b) Does strict 2PL lead to a deadlock?

## ii. Schedule 2:

```
R2(B), W2(B), R3(C), W3(C), R1(A), R1(B), W1(A), W1(B), Co1, R2(C), W2(C), Co2, R3(A), W3(A), Co3
```

- a) Is it possible under strict 2PL?
- b) Does strict 2PL lead to a deadlock?

- (b) A read-only transaction is a transaction that only reads from the database, without writing/inserting deleting. Answer the questions below by circling the correct answer.
- i. If all transactions are read-only, then every schedule is serializable.

TRUE or FALSE

ii. Only one transaction can hold a shared lock on the same item at any time.

TRUE or FALSE

iii. Only one transaction can hold an exclusive lock on the same item at any time.

TRUE or FALSE

(c) Given the following three transactions:

T1: R(A), W(B), I(D), R(C)

T2: R(B), R(D), W(C)

T3: R(D), R(C), R(D), W(A)

Assume that R(X) reads all tuples in table X, W(X) updates all tuples in X, and I(X) inserts one new tuple in X.

Does there exist a schedule of the above transactions that would result in a deadlock if executed under strict 2PL with **only exclusive table locks**? If so, write such a schedule with lock and unlock operations and indicate why the transactions are deadlocked. Otherwise write "No". Use L1(A) to refer to T1 locking table A, and U1(A) for unlocking.

L1(A); R1(A); L2(B); R2(B); L3(D), R3(D), L3(C); R3(D); W3(A)